GASPI: Global Address Space Programming Interface

Specification of a PGAS API for communication

Version 1.01

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1 Introduction to GASPI

1.1 Overview and Goals

GASPI stands for Global Address Space Programming Interface and is a Partitioned Global Address Space (PGAS) API. It aims at extreme scalability, high flexibility and failure tolerance for parallel computing environments. GASPI aims to initiate a paradigm shift from bulk-synchronous two-sided communication patterns towards an asynchronous communication and execution model. To that end GASPI leverages remote completion and one-sided RDMA driven communication in a Partitioned Global Address Space.

GASPI is neither a new language (like Chapel from Cray), nor an extension to a language (like Co-Array Fortran or UPC). Instead—very much in the spirit of MPI—it complements existing languages like C/C++ or Fortran with a PGAS API which enables the application to leverage the concept of the Partitioned Global Address Space. GASPI is not limited to a single memory model, but rather provides configurable RDMA PGAS memory segments. GASPI allows application developers to map the memory heterogeneity of a modern supercomputer node to these PGAS segments. As an example GASPI allows users to map the main memory of a GPGPU or Xeon Phi to a specific segment, to configure a GASPI segment per memory controller in a CC-NUMA system or to map non-volatile RAM to a specific segment. All these segments can directly read and write from/to each other - within the node and across all nodes. GASPI is failure tolerant in the sense that it provides timeout mechanisms for all non-local procedures, failure detection and the possibility to adapt to shrinking or growing node sets.

1.2 History

The GASPI specification originates from the PGAS API of the Fraunhofer ITWM (Fraunhofer Virtual Machine, FVM), which has been developed since 2005. Starting from 2007 this PGAS API has evolved into a robust commercial product (called GPI) which is used in the industry projects of the Fraunhofer ITWM. GPI offers a highly efficient and scalable programming model for Partitioned Global Address Spaces and has replaced MPI completely at Fraunhofer ITWM. In 2011 the partners of Fraunhofer ITWM, Fraunhofer SCAI, TUD, T-Systems SfR, DLR, KIT, FZJ, DWD and Scapos have initiated and launched the GASPI project to define a novel specification for a PGAS API (GASPI, based on GPI) and to make this novel GASPI specification a reliable, scalable and universal tool for the HPC community.

1.3 Design goals

GASPI has been designed with the following goals in mind:

- Extreme scalability.
• Efficient one sided asynchronous remote read/write operations based on remote completion.

• Multi-segment support to support e.g. heterogeneous systems and NUMA-pinning.

• Dynamic allocation of segments.

• Timeout mechanisms to allow failure tolerant programming.

• Asynchronous collective operations for groups of processes.

• Flexibility in the number of message queues, the queue sizes, atomic operations etc.

• A maximum freedom to implementors, where details are left to the implementation.

• A strong standard library which takes care of convenience procedures and cosmetics. The specification should be simple and solid.

2 GASPI terms and conventions

This section describes notational terms and conventions used throughout the GASPI document.

2.1 Naming Conventions

All procedures are named in accordance with the following convention. The procedures have gaspi_ as a prefix. The prefix is followed by the operation name.

2.2 Procedure specification

GASPI has adopted the procedure specification of MPI. Similar to the MPI standard, procedures in GASPI hence are first specified using a language independent notation. Immediately below this, the arguments of the procedure are given and marked as IN or OUT. The meanings of these are:

• the call uses but does not update an argument marked IN. For the C procedures these arguments are const-correct.

• the call may update an argument marked OUT.

Similar to MPI, in GASPI the passing of aliased procedure parameters results in undefined behavior.

Below the procedure arguments the ANSI C version of the function is shown, and below this, a version of the same function is shown for Fortran 2003. For the latter the corresponding definitions and derived types have to be include via
2.3 Semantic terms

The following semantic terms are used throughout the document:

**non-blocking** A procedure is non-blocking if the procedure may return before the operation completes.

```
Time   Operation   Call   Wait
```

**blocking** A procedure is blocking if the procedure only returns after the operation has completed.

```
Time   Operation   Call
```

**time-based blocking** A procedure is time-based blocking if the procedure may return after the operation completes or after a given timeout has been reached. A corresponding return value is used to distinguish between the two cases.

```
Time   Operation   Call
```

**local** A procedure is local if completion of the procedure depends only on the local executing Gaspi process.

**non-local** A procedure is non-local if completion of the operation may depend on the existence (and execution) of a remote Gaspi process.

**collective** A procedure is collective if all processes in a process group need to invoke the procedure. A collective call may or may not be synchronising.

**predefined** A predefined type is a datatype with a predefined constant name.

**timeout** A timeout is a mechanism required by procedures that might block (see blocking above). Timeout here is defined as the maximum time (in milliseconds) a called procedure will wait for outstanding communication from other processes. The special value 0 (defined as GASPI_TEST) indicates that the procedure will complete all local operations. The procedure subsequently returns the current status without waiting for data from other processes (non-blocking). On the other hand the special value -1 (defined as GASPI_BLOCK) instructs the procedure to wait indefinitely.
(blocking). A number greater than 0 indicates the maximum time the procedure will wait for data from other ranks (time-based blocking). The timeouts hence are soft: The timeout value \( n \) does not imply that the called procedure will return after \( n \) milliseconds. It just means that the procedure should wait for at most \( n \) milliseconds for data from other processes.

**synchronous** A procedure is called synchronous if progress towards completion only is achieved as long as the application is inside (executing) the procedure.

<table>
<thead>
<tr>
<th>Time</th>
<th>Progress</th>
<th>Calls</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

**asynchronous** A procedure is called asynchronous if progress towards completion may be achieved after the procedure exits.

<table>
<thead>
<tr>
<th>Time</th>
<th>Progress</th>
<th>Calls</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Please note that some of the semantic terms are not exclusive. Some of them do overlap. According to the definition, a collective procedure may also be a local procedure. Furthermore, a blocking procedure is per definition also a synchronous procedure; the reverse statement is not true.

## 2.4 Examples

The examples in this document are for illustration purposes only. They are not intended to specify the semantics.

## 3 GASPI concepts

### 3.1 Introduction and overview

In this section, the basic GASPI concepts are introduced. A more detailed description with the corresponding procedure specifications can be found in the subsequent topic-specific sections.

GASPI is a communication API that implements a Partitioned Global Address Space (PGAS) model. Each GASPI process may host parts (called segments) of the global address space. A local segment can be accessed with standard load/store operations and remote segments can be accessed by every thread of every GASPI process using the GASPI read and write operations.

GASPI was designed with remote direct memory access (RDMA) in mind. A network infrastructure that supports RDMA guarantees asynchronous and one-sided communication operations without involving the CPU. This is one of the
3.2 **GASPI processes**

**GASPI** preserves the concept of ranks. Each **GASPI** process receives a unique rank that identifies it during its runtime.

3.3 **GASPI groups**

A group is a subset of all processes. The group members have common collective operations. A collective operation is then restricted to the processes forming the group.

3.4 **GASPI segments**

Modern hardware typically involves a hierarchy of memory with respect to the bandwidth and latencies of read and write accesses. Within that hierarchy are non-uniform memory access (NUMA) partitions, solid state devices (SSDs), graphical processing unit (GPU) memory or many integrated cores (MIC) memory. The **GASPI** memory segments are supposed to map this variety of hardware layers to the software layer. In the spirit of the PGAS approach, these **GASPI** segments may be globally accessible from every thread of every **GASPI** process. **GASPI** segments can also be used to leverage different memory models within a single application or to even run different applications in a single Partitioned Global Address Space.

3.5 **GASPI one-sided communication**

One-sided asynchronous communication is the basic communication mechanism provided by **GASPI**. The one-sided communication comes in two flavors. There are read and write operations from and into the Partitioned Global Address Space. For the write operations **GASPI** makes use of the concept of remote completion in the form of so-called notifications. One-sided operations are non-blocking and asynchronous, allowing the program to continue its execution along the data transfer. The actual data transfer is managed by the underlying network infrastructure.

3.6 **GASPI queues**

**GASPI** offers the possibility to use different queues to handle the communication requests. The requests can be submitted to one of the supported queues. These queues allow more scalability and can be used as channels for different types of requests where similar types of requests are queued and then get synchronised together but independently from the other ones (separation of concerns).
specification guarantees fairness of transfers posted to different queues, i.e. no queue should see its communication requests delayed indefinitely.

Listing 1: Allgather with one-sided writes.

```
let nProc be the number of processes;
let iProc be the unique id of this process;
let src be the data to be distributed;
let dst be an array storing the destination addresses;

foreach process p in [0,nProc):
    write src into dst[p][iProc];
    // ^^^^^^ |
    // | remote address if p != iProc
    // wait for the completion of the writes;
barrier;
// the writes of all processes are completed
```

### 3.7 GASPI passive communication

Passive communication has a two-sided semantic, where there is a matching receive operation to a send request. Passive communication aims at communication patterns where the sender is unknown (i.e. it can be any process from the receiver perspective) but there is potentially the need for synchronisation between different processes.

The receive operation is a blocking call that has as low interference as possible (e.g. consumes no CPU cycles) and is ideally woken up by the network layer. This passive communication allows for fair distributed updates of globally shared parts of data.

### 3.8 GASPI global atomics

GASPI provides atomic operations for integral types, i.e. such variables can be manipulated atomically without fear of preemption causing corruption. There are two basic atomic operations: `fetch_and_add` and `compare_and_swap`. The values can be used as global shared variables and to synchronise processes or events.

The specification guarantees fairness, i.e. no process should see its atomic operation delayed indefinitely.

Listing 2: Dynamic work distribution: Clients atomically fetch a packet id and increment the value.

```
do
{
    packet := fetch_and_add (1);
```
3.9 Gaspi timeouts

Failure tolerant parallel programs necessitate non-blocking communication calls. Hence, Gaspi provides a timeout mechanism for all potentially blocking procedures.

Timeouts for procedures are specified in milliseconds. GASPI_BLOCK is a pre-defined timeout value which blocks the procedure call until completion. This value should not be used in failure tolerant programs, as it can block for an indefinitely amount of time in case of an error.

GASPI_TEST is another predefined timeout value which blocks the procedure for the shortest time possible, i.e. the time in which the procedure call processes an atomic portion of its work.

Examples:

Listing 3: Blocks until the communication queue is empty and may block indefinitely in case of a failure.
```
WAIT (... , GASPI_BLOCK);
```

Listing 4: Just check if the operation has completed and return as soon as possible.
```
WAIT (... , GASPI_TEST);
```

Listing 5: Blocks until the queue is empty or more than 10 milliseconds have passed since wait has been called.
```
WAIT (... , 10);
```

3.10 Gaspi collective communication

Collective communication is communication which involves a group of Gaspi processes. It is collective only for that group.

Collective operations can be either synchronous or asynchronous. Synchronous implies that progress is achieved only as long as the application is inside of the call. The call itself, however, may be interrupted by a timeout. The operation is then continued in the next call of the procedure. This implies that a collective operation may involve several procedure calls until completion.

Collective operations are exclusive per group, i.e. only one collective operation of a specific type can run at a given time for a given group. For example, two
allreduce for one group cannot run at the same time; however, an allreduce operation and a barrier can run at the same time.

Implementor advice: GASPI does not regulate whether individual collective operations should internally be handled synchronously or asynchronously, however: GASPI aims at an efficient, low-overhead programming model. If asynchronous operation is supported, it should leverage external network-resources, rather than consuming CPU cycles.

GASPI supports the following collective operations: barriers, reductions with predefined operations, reductions with user defined operations.

Collective operations have their own queue and hence typically will be synchronised independently from the operations on other queues (separation of concerns).

3.11 GASPI return values

GASPI procedures have three general return values:

GASPI_SUCCESS implies that the procedure has completed successfully.

GASPI_TIMEOUT implies that the procedure could not complete in the given period of time. This does not necessitate an error. The procedure has to be invoked subsequently in order to fully complete the operation.

GASPI_ERROR implies that the procedure has terminated due to an error. There are no predefined error values specifying the detailed cause of an error. gaspi_error_message translates the error code into a human readable format.

Implementor advice: An implementation may provide specific error values. All error codes in the range $[-1, \ldots, -999]$ are reserved and must not be used. If there are predefined error codes, each of the return codes must have a corresponding error message.

Additionally, each process has a state vector that contains the health state for all processes. The state vector is set after non-local operations and can be used to detect failures on remote processes.

4 GASPI definitions

4.1 Types

gaspi_rank_t

The GASPI rank type.
### 4.1 Types

<table>
<thead>
<tr>
<th>Type Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>gaspi_segment_id_t</td>
<td>The GASPI memory segment ID type.</td>
</tr>
<tr>
<td>gaspi_offset_t</td>
<td>The GASPI offset type. Offsets are measured relative to the beginning of a memory segment in units of bytes.</td>
</tr>
<tr>
<td>gaspi_size_t</td>
<td>The GASPI size type. Sizes are measured in units of bytes.</td>
</tr>
<tr>
<td>gaspi_queue_id_t</td>
<td>The GASPI queue ID type.</td>
</tr>
<tr>
<td>gaspi_notification_t</td>
<td>The GASPI notification type.</td>
</tr>
<tr>
<td></td>
<td>Implementor advice: The sum of the sizes of gaspi_notification_t and gaspi_tag_t should be at most 8 bytes in order to allow for Infiniband specific optimizations.</td>
</tr>
<tr>
<td>gaspi_notification_id_t</td>
<td>The GASPI notification ID type.</td>
</tr>
<tr>
<td></td>
<td>Implementor advice: The sum of the sizes of gaspi_notification_t should be at most 8 bytes in order to allow for Infiniband specific optimizations.</td>
</tr>
<tr>
<td>gaspi_atomic_value_t</td>
<td>The GASPI global atomic value type. An atomic value is unsigned and its maximum value can be queried using gaspi_atomic_max.</td>
</tr>
<tr>
<td>gaspi_return_t</td>
<td>The GASPI return value type.</td>
</tr>
<tr>
<td>vector&lt;gaspi_return_t&gt;</td>
<td>The vector type with return codes for individual processes. The length of the vector equals the number of processes in the GASPI program.</td>
</tr>
<tr>
<td>gaspi_timeout_t</td>
<td></td>
</tr>
</tbody>
</table>
4.2 Constants

The GASPI timeout type.

`gaspi_number_t`

A type that is used to count elements. That could be numbers of queues as well as the size of individual queues.

`gaspi_group_t`

The GASPI group type.

`gaspi_pointer_t`

A type that can point to some (area of) memory.

`gaspi_const_pointer_t`

A type that can point to some (area of) memory that cannot be modified using this pointer.

`gaspi_alloc_t`

The GASPI allocation policy type.

`gaspi_network_t`

The GASPI network infrastructure type.

`gaspi_string_t`

The GASPI constant string type.

`gaspi_statistic_counter_t`

The GASPI statistic counter type.

4.2 Constants

4.2.1 Timeout values

`GASPI_BLOCK`

`GASPI_BLOCK` is a timeout value which blocks a procedure call until completion.

`GASPI_TEST`

`GASPI_TEST` is a timeout value which blocks a procedure call for the shortest time possible.
4.2.2 Function return values

- **GASPI_SUCCESS**
  
  GASPI_SUCCESS is returned if a procedure call is completed successfully.

- **GASPI_TIMEOUT**
  
  GASPI_TIMEOUT is returned if a procedure call ran into a timeout.

- **GASPI_ERROR**
  
  GASPI_ERROR is returned if a procedure call finished with an error.

4.2.3 State vector states

- **GASPI_STATE_HEALTHY**
  
  GASPI_STATE_HEALTHY implies that a remote GASPI process is healthy and communication is possible.

- **GASPI_STATE_CORRUPT**
  
  GASPI_STATE_CORRUPT implies that the remote GASPI process is corrupted and communication is impossible.

4.2.4 Allocation policies

- **GASPI_ALLOC_DEFAULT**
  
  The GASPI_ALLOC_DEFAULT policy uses the operating systems default memory allocation policy.

  Implementor advice: A GASPI implementation is free to provide additional allocation policies.

4.2.5 Statistics interface

A GASPI implementation is free to define constants of the type gaspi_statistic_counter_t for specific statistics.

5 Execution model

5.1 Introduction and overview

GASPI allows both SPMD (Single Program, Multiple Data) and MPMD (Multiple Program, Multiple Data) style program execution. Hence, either a single
program or different programs can be started on the computational units. How a GAsPI application is started and initialized is implementation specific.

A rank is attributed to each created process. Ranks are a central aspect as they allow applications to identify processes and therefore allow to distribute work among the processes.

Furthermore, GAsPI provides segments. Segments are globally accessible memory regions. In general, the execution of a GAsPI process can be considered as split into several consecutive phases:

- **Setup (optional)**
  - Setting up configuration parameters
  - Performing environment checks

- **Initialization**
  - Initialization of the runtime environment
  - Creation of segments or groups (optional)

- **Working (optional)**
  - Communication calls
  - Collective operations
  - Atomic operations

- **Shutdown**
  - Cleanup of communication infrastructure

In the **setup** phase, the application may retrieve and modify the GAsPI configuration structure (see Sect. 5.2.1) determining the GAsPI runtime behavior. Optionally (but advisable), the application can perform environment checks (see Sect. 13) to make sure the application can be started safely and correctly.

In the **initialization** phase, the GAsPI runtime environment is set up in accordance with the parameters of the configuration structure by invocation of the initialization procedure. The initialization procedure is called before any other functionality, with the exception of pre-initialization routines for environment checking and declaration and retrieval of configuration parameters. After the initialization routine has been called, an optional step to perform is the creation of one or more segments and the creation of one or more groups. Segments are contiguous blocks of memory that may be accessed globally by all processes and where global data should be placed.

After the initialization, the application can proceed with its **working** phase and use the functionalities of GAsPI (communication, collectives, atomic operations, etc.).

The application should call the **shutdown** procedure (see Sect. 5.3.4) before it is terminated so that all resources and the communication infrastructure is cleaned up.
5.2 Process configuration

The entire set of execution phases define the GASPI life cycle. In principle, several life cycles can be invoked in one GASPI program.

Calling a routine in an execution phase in which it is not supposed to be executed results in undefined behavior.

5.2 Process configuration

5.2.1 GASPI configuration structure

The GASPI configuration structure describes the configuration parameters which influence the GASPI runtime behavior.

Please note, that for simplicity of notation this is a C-style definition. In bindings to other languages corresponding definitions will be used.

```
typedef struct {
    // maximum number of groups
gaspi_number_t group_max;

    // maximum number of segments
gaspi_number_t segment_max

    // one-sided comm parameter
gaspi_number_t queue_num;
gaspi_number_t queue_size_max;
gaspi_size_t transfer_size_max;

    // notification parameter
gaspi_number_t notification_num;

    // passive comm parameter
gaspi_number_t passive_queue_size_max;
gaspi_size_t passive_transfer_size_max;

    // collective comm parameter
gaspi_size_t allreduce_buf_size;
gaspi_number_t allreduce_elem_max;

    // network selection parameter
gaspi_network_t network;

    // communication infrastructure build up notification
gaspi_number_t build_infrastructure;

    void * user_defined;
} gaspi_config_t;
```

The definition of each of the configuration structure fields is as follows:
5.2 Process configuration

**group_max** the desired maximum number of permissible groups per process. There is a hardware/implementation dependent maximum.

**segment_max** the desired number of maximally permissible segments per GASPI process. There is a hardware/implementation dependent maximum.

**queue_num** the desired number of one-sided communication queues to be created. There is a hardware/implementation dependent maximum.

**queue_size_max** the desired number of simultaneously allowed on-going requests on a one-sided communication queue. There is a hardware/implementation dependent maximum.

**transfer_size_max** the desired maximum size of a single data transfer in the one-sided communication channel. There is a hardware/implementation dependent maximum.

**notification_num** the desired number of internal notification buffers for weak synchronisation to be created. There is a hardware/implementation dependent maximum.

**passive_queue_size_max** the desired number of simultaneously allowed on-going requests on the passive communication queue. There is a hardware/implementation dependent maximum.

**passive_transfer_size_max** the desired maximum size of a single data transfer in the passive communication channel. There is a hardware/implementation dependent maximum.

**allreduce_elem_max** the maximum number of elements in gaspi_allreduce. There is a hardware/implementation dependent maximum.

**allreduce_buf_size** the size of the internal buffer of gaspi_allreduce_user. There is a hardware/implementation dependent maximum.

**network** the network type to be used.

**build_infrastructure** indicates whether the communication infrastructure should be built up at startup time. The default value is true.

**user_defined** some user defined information that is application / implementation dependent.

The default configuration structure can be retrieved by gaspi_config_get. Its default values are implementation dependent. If some of the parameters are set by the program and assigned with gaspi_config_set, the requested values are just proposals. Depending on the underlying hardware capabilities, the implementation is allowed to overrule these proposals. gaspi_config_set has to be used in order to commit modifications of the configuration structure before the initialization routine is invoked. The actual values of the parameters can be retrieved by the corresponding GASPI getter routines (see Sect. 12) after the successful program initialization. The values of the configuration structure parameters need to be the same on all GASPI processes.
5.2 Process configuration

The user has the possibility to set the values on her own or leave the default values. Each field (where applicable) also has a maximum value to avoid user errors that might lead to too much instability or scalability problems (for example, the number of queues).

5.2.2 gaspi_config_get

The gaspi_config_get procedure is a synchronous local blocking procedure which retrieves the default configuration structure.

\[ \text{GASPI_CONFIG_GET ( config )} \]

**Parameter:**

(out) config: the default configuration

\[
gaspi_return_t
gaspi_config_get ( gaspi_config_t *config )
\]

\[
\text{function gaspi_config_get(config) 
& & result( res ) bind(C, name="gaspi_config_get")}
\text{type(gaspi_config_t) :: config}
\text{integer(gaspi_return_t) :: res}
\text{end function gaspi_config_get}
\]

**Execution phase:**

Setup

**Return values:**

GASPI_SUCCESS: operation has returned successfully
GASPI_ERROR: operation has finished with an error

After successful procedure completion, i.e. return value GASPI_SUCCESS, config represents the default configuration.

In case of error, the return value is GASPI_ERROR.

5.2.3 gaspi_config_set

The gaspi_config_set procedure is a synchronous local blocking procedure which sets the configuration structure for process initialization.

\[ \text{GASPI_CONFIG_SET ( config )} \]

**Parameter:**
(in) config: the configuration structure to be set

```
gaspi_return_t
gaspi_config_set ( gaspi_config_t const config )
```

```
function gaspi_config_set(new_config) &
& result( res ) bind(C, name="gaspi_config_set")
type(gaspi_config_t), value :: new_config
integer(gaspi_return_t) :: res
end function gaspi_config_set
```

**Execution phase:**

**Setup**

**Return values:**

GASPI_SUCCESS: operation has returned successfully

GASPI_ERROR: operation has finished with an error

After successful procedure completion, i.e. return value GASPI_SUCCESS, the runtime parameters for the GASPI process initialization are set in accordance with parameters of config.

In case of error, the return value is GASPI_ERROR.

### 5.3 Process management calls

#### 5.3.1 gaspi_proc_init

gaspi_proc_init implements the GASPI initialization of the application. It is a non-local synchronous time-based blocking procedure.

```
GASPI_PROC_INIT ( timeout )
```

**Parameter:**

(in) timeout: the timeout

```
gaspi_return_t
gaspi_proc_init ( gaspi_timeout_t const timeout )
```

```
function gaspi_proc_init(timeout_ms) &
& result( res ) bind(C, name="gaspi_proc_init")
integer(gaspi_timeout_t), value :: timeout_ms
integer(gaspi_return_t) :: res
end function gaspi_proc_init
```
5.3 Process management calls

Execution phase:
Initialization

Return values:
GASPI_SUCCESS: operation has returned successfully
GASPI_TIMEOUT: operation has run into a timeout
GASPI_ERROR: operation has finished with an error

The explicit start of a GASPI process or launch from command line is not specified. This is implementation dependent.

However, it is anticipated that gaspi_proc_init has information about the list of hosts on which the entire GASPI application is running either by environment variables, a command line argument, a daemon or some other mechanism. The actual transfer of knowledge is implementation dependent.

gaspi_proc_init registers a given process at the other remote GASPI processes and sets the corresponding entries in the state vector to a healthy state. If the parameter build_infrastructure in the configuration structure is set, also the communication infrastructure for passive and one-sided communication to all of the other processes is setup. Otherwise, there is no set up of communication infrastructure during the initialization. A rank is assigned to the given GASPI process in accordance with the position of the host in the list. The GASPI process running on the first host in the list has rank zero. The GASPI process running on the second host in the list has rank one and so on.

In case of a node failure, a GASPI process can be started on a new host, freshly allocated or selected from a set of pre-allocated spare hosts, by providing the list of machines in which the failed node is substituted by the new host. The new GASPI process then has the rank of the GASPI process which has been running on the failed node.

In case of the subsequent start of additional GASPI processes, the newly started GASPI process registers with the other remote GASPI processes. Note, that a subsequent change of the number of running GASPI processes invalidates GASPI_GROUP_ALL for the old running processes. Also the return value of gaspi_proc_num is changed.

The configuration structure should be created and modified by the application before calling the gaspi_proc_init procedure.

After successful procedure completion, gaspi_proc_init returns GASPI_SUCCESS and it guarantees that the application has been started on all hosts. In case that the build_infrastructure is set, return value GASPI_SUCCESS also implies that the communication infrastructure is up and ready to be used.

In case the application could not be initialized in line with the timeout parameter, the return value is GASPI_TIMEOUT. The application has not been initialized yet. A subsequent invocation is required to completely initialize the application.

In case of error, the return value is GASPI_ERROR. The application is not initialized.
Implementor advice: Calling `gaspi_proc_init` with an enabled parameter `build_infrastructure` is semantically equivalent to calling `gaspi_proc_init` with a disabled parameter `build_infrastructure` and subsequent calls to `gaspi_connect` in which an all-to-all connection is established.

User advice: For resource critical applications, it is recommended to disable the parameter `build_infrastructure` in the configuration structure.

User advice: A successful procedure completion does not mean that any communication or collective operation can already be used. Connections might need to be established. A segment has to be allocated for passive communication capabilities. If one-sided communication is supposed to be used, than the segment has to be registered in addition. If collective operations are needed, a group has to be created and committed.

### 5.3.2 `gaspi_proc_num`:

The total number of GASPI processes started, can be retrieved by `gaspi_proc_num`. This is a local synchronous blocking procedure.

```plaintext
GASPI_PROC_NUM ( proc_num )
```

**Parameter:**

`(out) proc_num`: the total number of GASPI processes.

```plaintext
gaspi_return_t
gaspi_proc_num ( gaspi_rank_t *proc_num )
```

```plaintext
function gaspi_proc_num(proc_num) &
& result( res ) bind(C, name="gaspi_proc_num")
  integer(gaspi_rank_t) :: proc_num
  integer(gaspi_return_t) :: res
end function gaspi_proc_num
```

**Execution phase:**

`Working`

**Return values:**

- `GASPI_SUCCESS`: operation has returned successfully
- `GASPI_ERROR`: operation has finished with an error

If successful, the return value is `GASPI_SUCCESS` and `gaspi_proc_num` retrieves the total number of processes that have been initialized and places this number in the `proc_num`.

In case of error, the return value is `GASPI_ERROR` and the value of `proc_num` is undefined.
5.3 Process management calls

5.3.3 gaspi_proc_rank

A rank identifies a GASPI process. The rank of a process lies in the interval \([0, P]\) where \(P\) can be retrieved through gaspi_proc_num. Each process has a unique rank associated with it. The rank of the invoking GASPI process can be retrieved by gaspi_proc_rank. It is a local synchronous blocking procedure.

\[
\text{GASPI\_PROC\_RANK ( rank )}
\]

Parameter:

(out) rank: the rank of the calling GASPI process.

\[
gaspi\_return\_t
\]
\[
gaspi\_proc\_rank ( gaspi\_rank\_t *rank )
\]

function gaspi_proc_rank(rank) &
& result(res) bind(C, name="gaspi_proc_rank")
integer(gaspi_rank_t) :: rank
integer(gaspi_return_t) :: res
end function gaspi_proc_rank

Execution phase:

Working

Return values:

GASPI\_SUCCESS: operation has returned successfully
GASPI\_ERROR: operation has finished with an error

gaspi_proc_rank retrieves, if successful, the rank of the calling process, placing it in the parameter rank and returning GASPI\_SUCCESS.

In case of error, the return value is GASPI\_ERROR and the value of the rank is undefined.

5.3.4 gaspi_proc_term

The shutdown procedure gaspi_proc_term is a synchronous non-local time-based blocking operation that releases resources and performs the required clean-up. There is no definition in the specification of a verification of a healthy global state (i.e. all processes terminated correctly).

After a shutdown call on a given GASPI process, it is undefined behavior if another GASPI process tries to use any non-local GASPI functionality involving that process.

\[
\text{GASPI\_PROC\_TERM ( timeout )}
\]
5.3 Process management calls

Parameter:

(in) timeout: the timeout

```fortran
gaspi_return_t
gaspi_proc_term ( gaspi_timeout_t timeout )
```

```fortran
function gaspi_proc_term(timeout_ms) &
    & result( res ) bind(C, name="gaspi_proc_term")
    integer(gaspi_timeout_t), value :: timeout_ms
    integer(gaspi_return_t) :: res
end function gaspi_proc_term
```

Execution phase:

Shutdown

Return values:

GASPI_SUCCESS: operation has returned successfully

GASPI_TIMEOUT: operation has run into a timeout

GASPI_ERROR: operation has finished with an error

In case of successful procedure completion, i.e., return value GASPI_SUCCESS, the allocated GASPI specific resources of the invoking GASPI process have been released. That means in particular that the communication infrastructure is shut down, all committed groups are released and all allocated segments are freed.

In case of timeout, i.e., return value GASPI_TIMEOUT, the local resources of the invoking GASPI process could not be completely released in the given period of time. A subsequent invocation is required to completely release all of the resources.

In case of error, i.e., return value GASPI_ERROR, the resources of the local GASPI process could not be released. The process is in an undefined state.

5.3.5 gaspi_proc_kill

`gaspi_proc_kill` sends an interrupt signal to a given GASPI process. It is a **synchronous non-local time-based blocking** procedure.

```fortran
GASPI_PROC_KILL ( rank
                   , timeout )
```

Parameter:

(in) rank: the rank of the process to be killed

(in) timeout: the timeout
5.3 Process management calls

```fortran
function gaspi_proc_kill(rank,timeout_ms) &
        result(res) bind(C, name="gaspi_proc_kill")
  integer(gaspi_rank_t), value :: rank
  integer(gaspi_timeout_t), value :: timeout_ms
  integer(gaspi_return_t) :: res
end function gaspi_proc_kill
```

Execution phase:

Working

Return values:

GASPI_SUCCESS: operation has returned successfully
GASPI_TIMEOUT: operation has run into a timeout
GASPI_ERROR: operation has finished with an error

`gaspi_proc_kill` sends an interrupt signal to the GASPI process incorporating the rank given by parameter `rank`. This can be used, for example, to realize the registration of a user-defined signal handler function which ensures the controlled shut down of an entire GASPI application at the global level if the application receives an interrupt signal (`STRG + C`) in the interactive master process. Every GASPI application should register such or a similar signal handler (c.f. listing 9).

In case of successful procedure completion, i.e. return value GASPI_SUCCESS, the remote GASPI process has been terminated.

In case of timeout, i.e. return value GASPI_TIMEOUT, the remote GASPI process could not be terminated in the given time. A subsequent invocation of the procedure is needed in order to complete the operation.

In case of error, i.e. return value GASPI_ERROR, the state of the remote GASPI process is undefined.

User advice: The kill signal terminates a GASPI process in an uncontrolled way. In this case, in order to provide a clean shutdown, it is advisable to register a user-defined signal callback function which guarantees a clean shutdown.

5.3.6 Example

The listing 7 shows a GASPI "Hello world" example. Please note that this example does not deal with failures.

Listing 7: GASPI hello world example.
5.3 Process management calls

```c
#include <stdio.h>
#include <stdlib.h>
#include <GASPI.h>

int main (int argc, char *argv[])
{
    gaspi_proc_init (GASPI_BLOCK);
    gaspi_rank_t iProc;
    gaspi_rank_t nProc;
    gaspi_proc_rank (&iProc);
    gaspi_proc_num (&nProc);
    printf ("Hello world from rank %i of %i!
", iProc, nProc);
    gaspi_proc_term (GASPI_BLOCK);
    return EXIT_SUCCESS;
}
```

Correspondingly the fortran version of GASPI "Hello world" assumes the form listing 8

```
program hello_world
use gaspi_c_binding
implicit none
integer(gaspi_return_t) :: res
integer(gaspi_rank_t) :: rank, num
integer(gaspi_timeout_t) :: timeout

timeout = GASPI_BLOCK

res = gaspi_proc_init(timeout)
res = gaspi_proc_rank(rank)
res = gaspi_proc_num(num)

print *,"Hello world from rank ",rank
res = gaspi_proc_term(timeout)
end program hello_world
```

The listing 9 shows the registration of a user defined signal handler function which ensures the controlled shut down of an entire GASPI application at the global level if the application receives an interrupt signal ($\text{STRG} + C$) in the...
interactive master process. Every GASPI application should register such or a similar signal handler.

Listing 9: Signal handling.
```
#include <signal.h>
#include <stdlib.h>
#include <GASPI.h>

void signalHandler (int sigint)
{
  gaspi_rank_t iProc;
  gaspi_rank_t nProc;

  gaspi_proc_rank (&iProc);
  gaspi_proc_num (&nProc);

  if (0 == iProc)
  {
    for (iProc = 1; iProc < nProc; ++iProc)
    {
      gaspi_proc_kill (iProc, GASPI_BLOCK);
    }
  }
  gaspi_proc_term (GASPI_BLOCK);
  exit (EXIT_FAILURE);
}

int main (int argc, char *argv[])
{
  gaspi_proc_init (GASPI_BLOCK);
  signal (SIGINT, &signalHandler);
  /* working phase */
  gaspi_proc_term (GASPI_BLOCK);
  return EXIT_SUCCESS;
}
```
5.4 Connection management utilities

5.4.1 gaspi_connect

In order to be able to communicate between two GASPI processes, the communication infrastructure has to be established. This is achieved with the synchronous non-local time-based blocking procedure gaspi_connect. It is bound to the working phase of the GASPI life cycle.

\[
\text{GASPI\_CONNECT ( rank, timeout )}
\]

Parameter:

(in) rank: the remote rank with which the communication infrastructure is established

(in) timeout: The timeout for the operation

\[
gaspi\_connect ( \text{gaspi\_rank\_t rank, gaspi\_timeout\_t timeout } )
\]

function gaspi_connect(rank,timeout_ms) &
& result( res ) bind(C, name="gaspi_connect")
integer(gaspi_rank_t), value :: rank
integer(gaspi_timeout_t), value :: timeout_ms
integer(gaspi_return_t) :: res

Execution phase:

Working

Return values:

GASPI\_SUCCESS: operation has returned successfully
GASPI\_TIMEOUT: operation has run into a timeout
GASPI\_ERROR: operation has finished with an error

gaspi_connect builds up the communication infrastructure, passive as well as one-sided and atomic operations, between the local and the remote GASPI process representing rank rank. The connection is bi-directional, i.e. it is sufficient that gaspi_connect is invoked by only one of the connection partners.

In case of successful procedure completion, i.e. return value GASPI\_SUCCESS, the communication infrastructure is established. If there is an allocated segment, the segment can be used as a destination for passive communication between the two nodes. In case the connection has already been established, e.g. by the connection partner, the return value is GASPI\_SUCCESS.
In case of return value `GASPI_TIMEOUT`, the communication infrastructure could not be established between the local GASPI process and the remote GASPI process in the given period of time.

In case of return value `GASPI_ERROR`, the communication infrastructure could not be established between the local GASPI process and the remote GASPI process.

In case of the latter two return values, a check of the state vector by invocation of `gaspi_state_vec_get` gives information on whether the remote GASPI process is still healthy.

**User advice:** Under the assumption that the GASPI process is initialized with parameter `build_infrastructure` set to `true`, all the connections are set up at initialization time. Hence, a subsequent call to `gaspi_connect` is superfluous in this case.

### 5.4.2 gaspi_disconnect

The `gaspi_disconnect` procedure is a synchronous local blocking procedure which disconnects a given process, identified by its rank, and frees all associated resources.

It is bound to the working phase of the GASPI life cycle.

```plaintext
GASPI_DISCONNECT ( rank , timeout )
```

**Parameter:**

*(in) rank:* the remote rank from which the communication infrastructure is disconnected

*(in) timeout:* The timeout for the operation

```plaintext
gaspi_disconnect ( gaspi_rank_t rank , gaspi_timeout_t timeout )
```

```plaintext
function gaspi_disconnect(rank,timeout_ms) & &
  result( res ) bind(C, name="gaspi_disconnect")
  integer(gaspi_rank_t), value :: rank
  integer(gaspi_timeout_t), value :: timeout_ms
  integer(gaspi_return_t) :: res
end function gaspi_disconnect
```

**Execution phase:**

Working

**Return values:**
GASPI_SUCCESS: operation has returned successfully
GASPI_TIMEOUT: operation has run into a timeout
GASPI_ERROR: operation has finished with an error

gaspi_disconnect disconnects the communication infrastructure, passive as well as one-sided and atomic operations, between the local and the remote GASPI process representing rank `rank`. The connection is bi-directional, i.e. it is sufficient if gaspi_disconnect is invoked by only one of the connection partners.

In case of successful procedure completion, i.e. return value GASPI_SUCCESS, the communication infrastructure is disconnected. Associated resources are freed on the local as well as on the remote side. In case the connection has already been disconnected, e.g. by the connection partner, the return value is GASPI_SUCCESS.

In case of error the return value is GASPI_ERROR.

In case of return value GASPI_TIMEOUT, the connection between the local GASPI process and the remote GASPI process could not be disconnected in the given period of time.

In case of the latter two return values local resources are freed and a check of the state vector by invocation of gaspi_state_vec_get gives information whether the remote GASPI process is still healthy.

After successful procedure completion, i.e. return value GASPI_SUCCESS, the connection is disconnected and can no longer be used.

5.5 State vector for individual processes

5.5.1 Introduction

A necessary pre-condition for realising a failure tolerant code is a detailed knowledge about the state of the communication partners of each local GASPI process.

GASPI provides a predefined type to describe the state of a remote GASPI process, which is the gaspi_state_t type. gaspi_state_t can have one of two values:

GASPI_STATE_HEALTHY implies that the remote GASPI process is healthy, i.e. communication is possible.

GASPI_STATE_CORRUPT means that the remote GASPI process is corrupted, i.e. there is no communication possible.

typedef vector<gaspi_state_t> gaspi_state_vector_t

gaspi_state_vector_t is a vector with state information for individual processes. The length of the vector equals the number of processes in the GASPI program and the entries are ordered based on the process ranks, i.e. entry 0 of the vector represents the state of process with the rank 0.
There are procedures to query the state of the communication partners after a
given communication request and also to reset the state after successful recovery.
These are described in the following subsections.

The state vector does not provide a global view, instead each process has its
own state vector that may be different to the state vector of another process.

5.5.2  gaspi_state_vec_get

The state vector is obtained by the *local synchronous blocking* function gaspi_
state_vec_get.

The state vector represents the states of all GASPI processes.

```
GASPI_STATE_VEC_GET ( state_vector )
```

**Parameter:**

*out* returns: the vector with individual return codes

```
gaspi_return_t gaspi_state_vec_get ( gaspi_state_vector_t *state_vector )
```

```
function gaspi_state_vec_get(state_vector) 
& 
    result( res ) bind(C, name="gaspi_state_vec_get")
    type(c_ptr), value :: state_vector
    integer(gaspi_return_t) :: res
end function gaspi_state_vec_get
```

**Execution phase:**

Working

**Return values:**

GASPI_SUCCESS: operation has returned successfully
GASPI_ERROR: operation has finished with an error

The state vector has one entry for each rank. It is created and initialized
during gaspi_proc_init. It is updated, case required, in each of the following
operations:

- group commitment
- gaspi_group_commit
- segment registration
- gaspi_segment_register
5.6 MPI Interoperability

Gaspi aims at providing interoperability with MPI in order to allow for incremental porting of such applications.

The startup of mixed MPI and Gaspi code is achieved by invoking `gaspi_proc_init` in an existing MPI program. This way, MPI takes care of distributing...
and starting the binary and GASPI just takes care of setting up its internal infrastructure.

GASPI and MPI communication should not occur at the same time, i.e. only the program layout given in Listing 10 is supported.

```
mpi_startup;
/* MPI part, no ongoing GASPI communication... */
/* ...finish all ongoing MPI communication */
mpi_barrier;
/* no ongoing MPI communication */
gaspi_proc_init;
while (!done) {
    /* GASPI part, no ongoing MPI communication... */
    /* ...finish all ongoing GASPI communication */
gaspi_barrier;
    /* MPI part, no ongoing GASPI communication... */
    /* ...finish all ongoing MPI communication */
    mpi_barrier;
}
/* MPI part, no ongoing GASPI communication */
gaspi_proc_term;
/* MPI part, no ongoing GASPI communication */
mpi_shutdown;
```

5.7 Argument checks and performance

GASPI aims at high performance and does not provide any argument checks at procedure invocation per default.

*Implementor advice:* The implementation should provide a specific library which includes argument checks. The GASPI procedures should include out of bounds checks, there.
6 Groups

6.1 Introduction

Groups are subsets of the total number of GASPI processes. The group members have common collective operations. Each GASPI process may participate in more than one group.

The use-cases are the collective operations provided in section 11 that make sense to be performed only for a subset of GASPI processes in order to avoid a complete (all processes) collective synchronisation point.

A group has to be defined and declared in each of the participating GASPI processes. Defining a group is a three step procedure. An empty group has to be created first. Then the participating GASPI processes, represented by their ranks, have to be attached. The group definition is a local operation. In order to activate the group, the group has to be committed by each of the participating GASPI processes. This is a collective operation for the group. Only after a successful group commit, can the group be used for collective operations.

The maximum number of groups allowed per GASPI process is restricted by the implementation. A user defined value can be set with gaspi_config_set before initialization (gaspi_proc_init).

In case one group desintegrates due to some failure, the group has to be re-established. If there is a new process replacing the failed one, the group has to be defined and declared on the newly started GASPI process(es). Re-establishment of the group is then achieved by recommitment of the group by the GASPI processes which were still 'alive' (functioning) and by the newly started GASPI process.

6.2 GASPI group generics

6.2.1 GASPI group type

Groups are specified with a special group type gaspi_group_t.

6.2.2 GASPI_GROUP_ALL

GASPI_GROUP_ALL is a predefined default group that corresponds to the whole set of GASPI processes. This is to be used for collective operations that work for the whole system.

```
gaspi_group_t GASPI_GROUP_ALL;
```

User advice: Note that GASPI_GROUP_ALL is a group definition like any other sub group. In order to be used, GASPI_GROUP_ALL also has to be committed by gaspi_group_commit.
6.3 Group creation

6.3.1 gaspi_group_create

The gaspi_group_create procedure is a synchronous local blocking procedure which creates an empty group.

\[
\text{GASPI\_GROUP\_CREATE ( group )}
\]

Parameter:

\text{(out) group: the created empty group}

\[
gaspi\_return\_t
gaspi\_group\_create ( gaspi\_group\_t \*group )
\]

\[
\begin{array}{ll}
\text{function gaspi_group_create(group) \\ & \& \text{result ( res ) bind(C, name="gaspi_group_create")} \\
& \text{integer(gaspi_group_t) :: group} \\
& \text{integer(gaspi_return_t) :: res} \\
end function gaspi_group_create
\end{array}
\]

Execution phase:

Working

Return values:

GASPI\_SUCCESS: operation has returned successfully
GASPI\_ERROR: operation has finished with an error

After successful procedure completion, i.e. return value GASPI\_SUCCESS, \textit{group} represents an empty group without any members.

In case of error, the return value is GASPI\_ERROR.

6.3.2 gaspi_group_add

The gaspi_group_add procedure is a synchronous local blocking procedure which adds a given rank to an existing group.

\[
\text{GASPI\_GROUP\_ADD ( group , rank )}
\]

Parameter:

\text{(inout) group: the group to which the rank is added}
\text{(in) rank: the rank to add to the group}
6.3 Group creation

gaspi_return_t
gaspi_group_add ( gaspi_group_t group, gaspi_rank_t rank )

function gaspi_group_add(group,rank) &
&   result(res) bind(C, name="gaspi_group_add")
   integer(gaspi_group_t), value :: group
   integer(gaspi_rank_t), value :: rank
   integer(gaspi_return_t) :: res
end function gaspi_group_add

Execution phase:
Working

Return values:
GASPI_SUCCESS: operation has returned successfully
GASPI_ERROR: operation has finished with an error

After successful procedure completion, i.e. return value GASPI_SUCCESS, the Gaspi process with rank is added to group. Whenever you add a rank the list of ranks is sorted in ascending order.

In case of error, the return value is GASPI_ERROR.

6.3.3 gaspi_group_commit

The gaspi_group_commit procedure is a synchronous collective time-based blocking procedure which establishes a group.

GASPI_GROUP_COMMIT ( group, timeout )

Parameter:
(in) group: the group to commit
(in) timeout: the timeout

gaspi_return_t
gaspi_group_commit ( gaspi_group_t group, gaspi_timeout_t timeout )

function gaspi_group_commit(group,timeout_ms) &
&   result(res) bind(C, name="gaspi_group_commit")
   integer(gaspi_group_t), value :: group
   integer(gaspi_timeout_t), value :: timeout_ms
   integer(gaspi_return_t) :: res
end function gaspi_group_commit
6.4 Group deletion

Execution phase:
Working

Return values:
GASPI_SUCCESS: operation has returned successfully
GASPI_TIMEOUT: operation has run into a timeout
GASPI_ERROR: operation has finished with an error

The group committed by all participating processes must contain all ranks and
must identical for all processes, otherwise the result is undefined.

After successful procedure completion, i.e. return value GASPI_SUCCESS, the
group given by the parameter group is established. Collective operations invoked
by the members of the group are allowed from this moment on.

In case of timeout, i.e. return value GASPI_TIMEOUT, the group could not be
established on all ranks forming the group in the given period of time. The
group is in an undefined state and collective operations on the group yield
undefined behavior. A subsequent invocation is required in order to completely
establish the group.

In case of error, i.e. return value GASPI_ERROR, the group could not be estab-
lished. The group is in an undefined state and collective operations defined on
the given group yield undefined behavior.

In both cases, GASPI_TIMEOUT and GASPI_ERROR, the Gaspi state vector should
be checked in order to eliminate the possibility of a failure.

User advice: Any group commit should be performed only by a single
thread of a process. If two Gaspi processes are members of two
groups, then the order of the group commits should be the same on
both processes in order to avoid deadlocks.

Implementor advice: If the parameter build_infrastructure is not set,
the procedure gaspi_group_commit must set up the infrastructure for
all possible operations of the group.

6.4 Group deletion

6.4.1 gaspi_group_delete

The gaspi_group_delete procedure is a synchronous local blocking procedure
which deletes a given group.

GASPI_GROUP_DELETE ( group )

Parameter:
(in) group: the group to be deleted
6.5 Group utilities

6.5.1 \texttt{gaspi\_group\_num}

The \texttt{gaspi\_group\_num} procedure is a \textit{synchronous local blocking} procedure which returns the current number of allocated groups.

\begin{verbatim}
gaspi\_return\_t
gaspi\_group\_delete ( gaspi\_group\_t group )
\end{verbatim}

\begin{verbatim}
function gaspi\_group\_delete(group) &
    & result( res ) bind(C, name="gaspi\_group\_delete")
    integer(gaspi\_group\_t), value :: group
    integer(gaspi\_return\_t) :: res
end function gaspi\_group\_delete
\end{verbatim}

\textit{Execution phase:}

\textit{Working}

\textit{Return values:}

\texttt{GASPI\_SUCCESS}: operation has returned successfully

\texttt{GASPI\_ERROR}: operation has finished with an error

After successful procedure completion, i.e. return value \texttt{GASPI\_SUCCESS}, \texttt{group} is deleted and cannot be used further.

In case of error, the return value is \texttt{GASPI\_ERROR}.

\textit{Implementor advice}: If the parameter \texttt{build\_infrastructure} is not set to true, the procedure \texttt{gaspi\_group\_delete} must disconnect all connections which have been set up in the call to \texttt{gaspi\_group\_commit} and free all associated resources.

6.5 Group utilities
6.5 Group utilities

Execution phase:
Working

Return values:
GASPI_SUCCESS: operation has returned successfully
GASPI_ERROR: operation has finished with an error

After successful procedure completion, i.e. return value GASPI_SUCCESS, group_num contains the current number of allocated groups. The value of group_num is related to the parameter group_max in the configuration structure and cannot exceed that value. The value can be implementation specific.

6.5.2 gaspi_group_size

The gaspi_group_size procedure is a synchronous local blocking procedure which returns the number of ranks of a given group.

GASPI_GROUP_SIZE ( group
, group_size )

Parameter:
(in) group: the group to be examined
(out) group_size: the number of ranks in a given group

gaspi_return_t
gaspi_group_size ( gaspi_group_t group
, gaspi_number_t *group_size )

function gaspi_group_size(group,group_size) &
&    result( res ) bind(C, name="gaspi_group_size")
    integer(gaspi_group_t), value :: group
    integer(gaspi_number_t) :: group_size
    integer(gaspi_return_t) :: res
end function gaspi_group_size

Execution phase:
Working

Return values:
GASPI_SUCCESS: operation has returned successfully
GASPI_ERROR: operation has finished with an error
After successful procedure completion, i.e. return value \texttt{GASPI\_SUCCESS}, the parameter \texttt{group\_size} contains the number of \texttt{GASPI} processes forming the \texttt{group}.

In case of error, the return value is \texttt{GASPI\_ERROR}. The parameter \texttt{group\_size} has an undefined value.

### 6.5.3 \texttt{gaspi\_group\_ranks}

The \texttt{gaspi\_group\_ranks} procedure is a \emph{synchronous local blocking} procedure which returns a list of ranks of \texttt{GASPI} processes forming the group.

\begin{verbatim}
GASPI\_GROUP\_RANKS ( group
                            , group_ranks[group\_size] )
\end{verbatim}

**Parameter:**

\begin{itemize}
  \item \texttt{(in) group:} the group to be examined
  \item \texttt{(out) group\_ranks:} the list of ranks forming the group
\end{itemize}

\begin{verbatim}
gaspi\_return\_t
gaspi\_group\_ranks ( gaspi\_group\_t group
                            , gaspi\_rank\_t \*group\_ranks )
\end{verbatim}

**Execution phase:**

Working

**Return values:**

\begin{itemize}
  \item \texttt{GASPI\_SUCCESS:} operation has returned successfully
  \item \texttt{GASPI\_ERROR:} operation has finished with an error
\end{itemize}

After successful procedure completion, i.e. return value \texttt{GASPI\_SUCCESS}, the list \texttt{group\_ranks} contains the ranks of the processes that belong to the \texttt{group}. The list is not allocated by the procedure. The list allocation is supposed to be done outside of the procedure. The size of the list can be inquired by \texttt{gaspi\_group\_size}.

In case of error, the return value is \texttt{GASPI\_ERROR}. The list \texttt{group\_ranks} has an undefined value.
7 GASPI segments

7.1 Introduction and overview

Modern hardware has a complex memory hierarchy with different bandwidth and latencies for read and write accesses. Among them are non-uniform memory access (NUMA) partitions, solid state devices (SSDs), graphical processing unit (GPU) memory or many integrated cores (MIC) memory.

The GASPI memory segments are thus an abstraction representing any kind of memory level, mapping the variety of hardware layers to the software layer. A segment is a contiguous block of virtual memory. In the spirit of the PGAS approach, these GASPI segments may be globally accessible from every thread of every GASPI process and represent the partitions of the global address space.

By means of the GASPI memory segments it is also possible for multiple memory models or indeed multiple applications to share a single Partitioned Global Address Space.

Since segment allocation is expensive and the total number of supported segments is limited due to hardware constraints, the GASPI memory management paradigm is the following. GASPI provides only a few relatively large segments. Allocations inside of the pre-allocated segment memory are managed by the application.

Every GASPI process may possess a certain number of segments (not necessarily equal to the number possessed by the other ranks) that may be accessed as common memory, whether locally—with normal memory operations—or remotely—with the communication routines of GASPI.

In order to use a segment for communication between two processes, some setup steps are required in general.

A memory segment has to be allocated in each of the processes by the local procedure `gaspi_segment_alloc`. In order to also use the segments for one-sided communication, the memory segment has to be registered on the remote process which will access the memory segment at some point. This is achieved by the non-local procedure `gaspi_segment_register`.

*User advice:* If the parameter `build_infrasturcture` is not set, a connection has to be established between the processes before the segment can be registered at the remote process. This is accomplished by calling the procedure `gaspi_connect`.

`gaspi_segment_create` unites these steps into a single collective procedure for an entire group. After successful procedure completion, a common segment is created on each GASPI process forming the group which can be immediately used for communication among the group members.

During the lifetime of an application no segment is available unless it is explicitly created with `gaspi_segment_alloc` or `gaspi_segment_create` after the GASPI startup.
## 7.2 Segment creation

### 7.2.1 gaspi_segment_alloc

The *synchronous local blocking* procedure `gaspi_segment_alloc` allocates a memory segment and optionally maps it in accordance with a given allocation policy.

```c
GASPI_SEGMENT_ALLOC ( segment_id
    , size
    , alloc_policy )
```

**Parameter:**

*(in) segment_id:* The segment ID to be created. The segment IDs need to be unique on each GASPI process.

*(in) size:* The size of the segment in bytes

*(in) alloc_policy:* allocation policy

```c
gaspi_return_t
gaspi_segment_alloc ( gaspi_segment_id_t segment_id
    , gaspi_size_t size
    , gaspi_alloc_t alloc_policy )
```

**Execution phase:**

**Working**

**Return values:**

GASPI_SUCCESS: operation has returned successfully

GASPI_ERROR: operation has finished with an error

`gaspi_segment_alloc` allocates a segment of size `size` that will be referenced by the `segment_id` identifier. This identifier parameter has to be unique in the local GASPI process. Creating a new segment with an existing segment ID results in undefined behavior. Note that the total number of segments is restricted by the underlying hardware capabilities. The maximum number of supported segments can be retrieved by invoking `gaspi_segment_max`. 
7.2 Segment creation

Allocation of segments in GASPI allows for various so-called policies. The default policy in a cc-numa mode for example might be an allocation of socket-local memory, a different policy might allow to map GPU memory into the main memory of the host and yet another policy might allow for a direct access of external non-volatile RAM.

The alloc_policy is used to pass an allocation policy. The default allocation policy behavior is left to the implementation. The default allocation parameter is GASPI_ALLOC_DEFAULT.

After successful procedure completion, i.e. return value GASPI_SUCCESS, the segment can be accessed locally. In case that there is a connection established to a remote GASPI process, it can also be used for passive communication between the two GASPI processes. (Note that this is always the case if the process has been initialized with the parameter build_infrastructure set to true), it can also be used for passive communication between the two GASPI processes; either as a source segment for gaspi_passive_send or as a destination segment for gaspi_passive_receive.

A return value GASPI_ERROR indicates that the segment allocation failed and the segment cannot be used.

User advice: A GASPI implementation may allocate more memory than requested by the application for internal management.

Implementor advice: In case of non-uniform memory access architectures, the memory should be allocated close to the calling process. The allocation policy of the calling process should not be modified.

7.2.2 gaspi_segment_register

In order to be used in a one-sided communication request on an existing connection, a segment allocated by gaspi_segment_alloc needs to be made visible and accessible for the other GASPI processes. This is accomplished by the procedure gaspi_segment_register. It is a synchronous non-local time-based blocking procedure.

\[
\text{GASPI\_SEGMENT\_REGISTER ( segment_id, rank, timeout )}
\]

Parameter:

(in) segment_id: The segment ID to be registered. The segment ID’s need to be unique for each GASPI process

(in) rank: The rank of the GASPI process which should register the new segment

(in) timeout: The timeout for the operation
7.2 Segment creation

\begin{verbatim}
gaspi_return_t
gaspi_segment_register ( gaspi_segment_id_t segment_id,
                        gaspi_rank_t rank,
                        gaspi_timeout_t timeout )

function gaspi_segment_register(segment_id,rank,timeout_ms) &
     & result( res ) bind(C, name="gaspi_segment_register")
     integer(gaspi_segment_id_t), value :: segment_id
     integer(gaspi_rank_t), value :: rank
     integer(gaspi_timeout_t), value :: timeout_ms
     integer(gaspi_return_t) :: res
end function gaspi_segment_register
\end{verbatim}

Execution phase:

Working

Return values:

GASPI_SUCCESS: operation has returned successfully
GASPI_TIMEOUT: operation has run into a timeout
GASPI_ERROR: operation has finished with an error

\begin{verbatim}
gaspi_segment_register makes the segment referenced by the segment_id identifier visible and accessible to the Gaspi process with the associated rank.
\end{verbatim}

User advice: If the parameter build_infrastructure is not set, a connection has to be established between the processes before the segment can be registered at the remote process. This is accomplished calling the procedure gaspi_connect.

In case of successful procedure completion, i.e. return value GASPI_SUCCESS, the local segment can be used for one-sided communication requests which are invoked by the given remote process.

In case of return value GASPI_TIMEOUT, the segment could not be registered in the given period of time. The segment cannot be used for one-sided communication requests which are invoked by the given remote process. A subsequent call of gaspi_segment_register has to be invoked in order to complete the registration request.

In case of return value GASPI_ERROR, the segment could not be registered on the remote side. The segment cannot be used for one-sided communication requests which are invoked by the given remote process.

In case of the latter two return values, a check of the state vector by invocation of gaspi_state_vec_get gives information as to whether or not the remote Gaspi process is still healthy.
User advice: Note that a local return value GASPI_SUCCESS does not imply that the remote process is informed explicitly that the segment is accessible. This can be achieved through an explicit synchronisation, either by one of the collective operations or by an explicit notification.

7.2.3 gaspi_segment_create

gaspi_segment_create is a synchronous collective time-based blocking procedure. It is semantically equivalent to a collective aggregation of gaspi_segment_alloc, gaspi_segment_register and gaspi_barrier involving all of the members of a given group. If the communication infrastructure was not established for all group members beforehand, gaspi_segment_create will accomplish this as well.

GASPI_SEGMENT_CREATE ( segment_id , size , group , timeout , alloc_policy )

Parameter:

(in) segment_id: The ID for the segment to be created. The segment ID’s need to be unique for each GASPI process

(in) size: The size of the segment in bytes

(in) group: The group which should create the segment

(in) timeout: The timeout for the operation

(in) alloc_policy: allocation policy

gaspi_return_t

gaspi_segment_create ( gaspi_segment_id_t segment_id , gaspi_size_t size , gaspi_group_t group , gaspi_timeout_t timeout , gaspi_alloc_t alloc_policy )

function gaspi_segment_create(segment_id,size,group, &
   & timeout_ms,alloc_policy) &
   & result ( res ) bind(C, name="gaspi_segment_create")
   integer(gaspi_segment_id_t), value :: segment_id
   integer(gaspi_size_t), value :: size
   integer(gaspi_group_t), value :: group
   integer(gaspi_timeout_t), value :: timeout_ms
   integer(gaspi_alloc_t), value :: alloc_policy
   integer(gaspi_return_t) :: res
end function gaspi_segment_create
7.2 Segment creation

Execution phase:
Working

Return values:
GASPI_SUCCESS: operation has returned successfully
GASPI_TIMEOUT: operation has run into a timeout
GASPI_ERROR: operation has finished with an error

gaspi_segment_create allocates a segment of size \textit{size} that will be referenced by the \textit{segment_id} identifier. This identifier parameter has to be unique on the local GASPI process. Creating a new segment with an existing segment ID results in undefined behavior. \texttt{gaspi_segment_create} makes the segment referenced by the \textit{segment_id} identifier visible and accessible to all of the GASPI processes forming the group \textit{group}. The maximum number of supported segments can be retrieved by invoking \texttt{gaspi_segment_max}. The \textit{alloc_policy} is used to pass an allocation policy. The default allocation policy behavior is left to the implementation.

After successful procedure completion, i.e. GASPI_SUCCESS, the segment can be accessed locally and it can be used as a destination for the passive communication channel. Either as a source segment for \texttt{gaspi_passive_send} or as a destination segment for \texttt{gaspi_passive_receive}. Furthermore, it can be used for one-sided communication requests, which are invoked by the remote processes forming the group \textit{group} or global atomic operations. The segment \textit{segment_id} is ready to be used.

For consistency and programs with hard failure tolerance requirements, the operation must be performed within \textit{timeout} milliseconds. In case of return value GASPI_TIMEOUT, progress has been achieved, however the operation could not be completed in the given timeout. The segment cannot be used locally neither remotely. The segment cannot be used for one-sided or passive communication requests which are invoked by the other remote processes forming the group. The same applies to global atomic operations. A subsequent call of \texttt{gaspi_segment_create} has to be invoked in order to complete the segment creation.

In case of return value GASPI_ERROR, the segment creation failed in one of the above progress steps on at least one of the involved GASPI processes. The segment cannot be used locally neither remotely. The segment cannot be used for one-sided or passive communication requests which are invoked by the other remote processes forming the group. The same applies to global atomic operations.

In case of the latter two return values, a check of the state vector by invocation of \texttt{gaspi_state_vec_get} gives information whether the involved remote GASPI processes are still healthy.

User advice: A GASPI implementation may allocate more memory than requested by the application for internal management.
Implementor advice: In case of non-uniform memory access architectures, the memory should be allocated close to the calling process. The allocation policy of the calling process should not be modified.

7.3 Segment deletion

7.3.1 gaspi_segment_delete

The synchronous local blocking procedure gaspi_segment_delete releases the resources of a previously allocated memory segment.

```
GASPI_SEGMENT_DELETE ( segment_id )
```

Parameter:

(in) `segment_id`: The segment ID to be deleted.

```
gaspi_return_t
gaspi_segment_delete ( gaspi_segment_id_t segment_id )
```

```
function gaspi_segment_delete(segment_id) &
     result(res) bind(C, name="gaspi_segment_delete")
     integer(gaspi_segment_id_t), value:: segment_id
     integer(gaspi_return_t):: res
end function gaspi_segment_delete
```

Execution phase:

Working

Return values:

GASPI_SUCCESS: operation has returned successfully
GASPI_ERROR: operation has finished with an error

gaspi_segment_delete releases the resources of the segment which is referenced by the `segment_id` identifier.

After successful procedure completion, i.e. return value GASPI_SUCCESS, the segment is deleted and the resources are released. It would be an application error to use the segment for communication between two GASPI processes after gaspi_delete has been called.

In case of return value GASPI_ERROR, the segment deletion failed. The segment is in an undefined state and cannot be used locally neither remotely. The segment cannot be used for one-sided or passive communication requests which are invoked by the other remote processes forming the group. The same applies to global atomic operations.
7.4 Segment utilities

7.4.1 gaspi_segment_num

The `gaspi_segment_num` procedure is a synchronous local blocking procedure which returns the current number of allocated segments.

\[ \text{GASPI\_SEGMENT\_NUM} ( \ segment\_num ) \]

Parameter:

\((\text{out})\ segment\_num:\) the current number of allocated segments

\[ \text{gaspi\_return\_t} \]
\[ \text{gaspi\_segment\_num} ( \ \text{gaspi\_number\_t} * \text{segment\_num} ) \]

function gaspi_segment_num(segment_num) &
& result( res ) bind(C, name="gaspi_segment_num")
  integer(gaspi_number_t) :: segment_num
  integer(gaspi_return_t) :: res
end function gaspi_segment_num

Execution phase:

Working

Return values:

GASPI\_SUCCESS: operation has returned successfully
GASPI\_ERROR: operation has finished with an error

After successful procedure completion, i.e. return value GASPI\_SUCCESS, segment_num contains the current number of locally allocated segments provided by Gaspi. The value of segment_num is related to the parameter segment_max in the configuration structure which is retrieved by gaspi\_config\_get and cannot exceed that value. The maximum number of allocatable segments per process might be implementation specific.

In case of error, the return value is GASPI\_ERROR. The parameter segment_num has an undefined value.

7.4.2 gaspi_segment_list

The `gaspi_segment_list` procedure is a synchronous local blocking procedure which returns a list of locally allocated segment IDs.

\[ \text{GASPI\_SEGMENT\_LIST} ( \ num \]
\[ , \ segment\_id\_list[num] \ ) \]
7.4 Segment utilities

Parameter:

*(in)* `num`: number of segment IDs to collect

*(out)* `segment_list[num]`: list of locally allocated segment IDs

```fortran
function gaspi_segment_list(num,segment_id_list) &
& result(res) bind(C, name="gaspi_segment_list")
  integer(gaspi_number_t), value :: num
  type(c_ptr), value :: segment_id_list
  integer(gaspi_return_t) :: res
end function gaspi_segment_list
```

Execution phase:

Working

Return values:

GASPI_SUCCESS: operation has returned successfully

GASPI_ERROR: operation has finished with an error

After successful procedure completion, i.e. return value GASPI_SUCCESS, `segment_list[num]` contains the IDs of `num` locally allocated segments.

The size of `segment_list[num]` needs to be at least `num` elements long.

In case of error, the return value is GASPI_ERROR. The parameter `segment_list[num]` has an undefined value.

7.4.3 gaspi_segment_ptr

Segments are identified by a unique ID. This ID can be used to obtain the virtual address of that local segment of memory. The procedure `gaspi_segment_ptr` returns the pointer to the segment represented by a given segment ID. It is a synchronous local blocking procedure.

```fortran
GASPI_SEGMENT_PTR ( segment_id
  , pointer )
```

Parameter:

*(in)* `segment_id`: The segment ID.

*(out)* `pointer`: The pointer to the memory segment.
gaspi_return_t
gaspi_segment_ptr ( gaspi_segment_id_t segment_id,
gasi_pointer_t *pointer )

function gaspi_segment_ptr(segment_id,ptr) &
& result( res ) bind(C, name="gaspi_segment_ptr")
  integer(gaspi_segment_id_t), value :: segment_id
  type(c_ptr) :: ptr
  integer(gaspi_return_t) :: res
end function gaspi_segment_ptr

Execution phase:

Working

Return values:

GASPI_SUCCESS: operation has returned successfully
GASPI_ERROR: operation has finished with an error

After successful procedure completion, i.e. GASPI_SUCCESS, the output parameter pointer contains the virtual address pointer of the memory identified by segment_id. This gaspi_pointer_t can then be used to reference the segment and perform memory operations.

In case of return value GASPI_ERROR, the translation of the segment ID to a pointer to a virtual memory address failed. The pointer contains an undefined value and cannot be used to reference the segment.

7.5 Segment memory management

Each thread of a process may have global read or write access to all of the segments provided by remote GASPI processes if there is a connection established between the processes and if the respective segments have been registered on the local process.

Since a segment is an entire contiguous block of virtual memory, allocations inside of the pre-allocated segment memory need to be managed.

GASPI does not provide dedicated memory management functionality for the local segments. This is left to the application. Since a default implementation for memory management cannot include knowledge about the specific problem, a good problem-related implementation of a memory management will always better than any predefined implementation.

Local and non-local GASPI procedures specify in general memory addresses within the Partitioned Global Address Space by the triple consisting of a rank, a segment identifier and an offset. This prevents a global all-to-all distribution of memory addresses, since memory addresses of memory segments could be and normally are different on different GASPI processes.
A local buffer is specified by the pair \textit{segment\_id, offset}. The buffer is located at address

\[
\text{buffer\_address} = \text{base\_addr(\text{segment\_id}) + offset}
\]

where \text{base\_addr(\text{segment\_id})} is the base address of the segment with identifier \text{segment\_id}. It can be obtained by applying \text{gaspi\_segment\_ptr} on the local process.

A remote buffer is specified by the triple \textit{remote\_rank, remote\_segment\_id, remote\_offset}. The address of the remote buffer can be calculated analogously to the local buffer. The only difference is the determination of the base address. Here, it is the address which would be obtained by invoking \text{gaspi\_segment\_ptr} on the remote \text{GASPI} process with \text{remote\_segment\_id} as input parameter.

8 One-sided communication

8.1 Introduction and overview

One-sided asynchronous communication is the basic communication mechanism provided by \text{GASPI}. Hereby, one \text{GASPI} process specifies all communication parameters, both for the local and the remote side. Due to the asynchronicity, a complete communication involves two procedure calls. First, one call to initiate the communication. This call posts a communication request to the underlying network infrastructure. The second call waits for the completion of the communication request.

For one-sided communication, \text{GASPI} provides the concept of communication queues. All operations placed on a certain queue \textit{q} by one or several threads are finished after a single wait call on the queue \textit{q} has returned successfully. Separation of concerns is possible by using different queues for different tasks, e.g. one queue for operations on data and another queue for operations on metadata.

The different communication queues guarantee fair communication, i.e. no queue should see its communication requests delayed indefinitely.

One-sided communication calls can basically be divided into two operation types: read and write. The read operations transfer data from a remote segment to a local segment. The write operations transfer data from a local segment to a remote segment.

The number of communication queues and their size can be configured at initialization time, otherwise default values will be used. The default values are implementation dependent. Maximum values are also defined.

For the write operation there are four different variants that allow different communication patterns:

- \text{gaspi\_write}
- \text{gaspi\_write\_notify}
8.2 Basic communication calls

8.2.1 gaspi_write

The simplest form of a write operation is gaspi_write which is a single communication call to write data to a remote location. It is an asynchronous non-local time-based blocking procedure.

```
GASPI_WRITE ( segment_id_local, offset_local, rank, segment_id_remote, offset_remote, size, queue, timeout )
```

Parameter:

- (in) segment_id_local: the local segment ID to read from
- (in) offset_local: the local offset in bytes to read from
- (in) rank: the remote rank to write to
- (in) segment_id_remote: the remote segment to write to
- (in) offset_remote: the remote offset to write to
- (in) size: the size of the data to write
8.2 Basic communication calls

**(in) queue:** the queue to use

**(in) timeout:** the timeout

```fortran
function gaspi_write(segment_id_local, offset_local, 
&    rank, segment_id_remote, offset_remote, size, 
&    queue, timeout_ms) & 
&    result ( res ) bind(C, name="gaspi_write")
  integer(gaspi_segment_id_t), value :: segment_id_local
  integer(gaspi_offset_t), value :: offset_local
  integer(gaspi_rank_t), value :: rank
  integer(gaspi_segment_id_t), value :: segment_id_remote
  integer(gaspi_offset_t), value :: offset_remote
  integer(gaspi_size_t), value :: size
  integer(gaspi_queue_id_t), value :: queue
  integer(gaspi_timeout_t), value :: timeout_ms
  integer(gaspi_return_t) :: res
end function gaspi_write
```

**Execution phase:**

Working

**Return values:**

GASPI_SUCCESS: operation has returned successfully
GASPI_TIMEOUT: operation has run into a timeout
GASPI_ERROR: operation has finished with an error

`gaspi_write` posts a communication request which asynchronously transfers a contiguous block of `size` bytes from a source location of the local Gaspi process to a target location of a remote Gaspi process. This communication request is posted to the communication queue `queue`. The source location is specified by the pair `segment_id_local, offset_local`. The target location is specified by the triple `rank, segment_id_remote, offset_remote`.

A valid `gaspi_write` communication request requires that the local and the remote segment are allocated, that there is a connection between the local and the remote Gaspi process and that the remote segment has been registered on
the local GASPI process. Otherwise, the communication request is invalid and the procedure returns with \texttt{GASPI\_ERROR}.

After successful procedure completion, i.e. return value \texttt{GASPI\_SUCCESS}, the communication request has been posted to the underlying network infrastructure. One new entry is inserted into the given queue.

Successive \texttt{gaspi\_write} calls posted to the same queue and the same destination rank are not guaranteed to be non-overtaking. However, a subsequent \texttt{gaspi\_notify}, which is posted to the same queue is guaranteed to be non-overtaking. In particular, one can hence assume, that if the corresponding notification has arrived on the remote process, the data from the earlier posted request to the same process has also arrived on the remote side.

\texttt{gaspi\_write} calls may be posted from every thread of the GASPI process.

If the procedure returns with \texttt{GASPI\_TIMEOUT}, the communication request could not be posted to the hardware during the given timeout. This can happen, if another thread is in a \texttt{gaspi\_wait} for the same queue. A subsequent call of \texttt{gaspi\_write} has to be invoked in order to complete the write call.

A communication request posted to a given queue can be considered as completed, if the correspondent \texttt{gaspi\_wait} returns with \texttt{GASPI\_SUCCESS}.

If the queue to which the communication request is posted is full, i.e. if the number of posted communication requests has already reached the queue size of a given queue, the communication request fails and the procedure returns with return value \texttt{GASPI\_ERROR}. If a saturated queue is detected, there are the following two options: Either one invokes a \texttt{gaspi\_wait} on the given queue in order to wait for all the posted requests to be finished. Alternatively it is possible to use another queue.

\textit{User advice:} Return value \texttt{GASPI\_SUCCESS} does not mean, that the data has been transferred or buffered or that the data has arrived at the remote side.

It is allowed to write data to the source location while the communication is ongoing. However, the result on the remote side would be some undefined interleaving of the data that was present when the call was issued and the data that was written later.

It is also allowed to read from the source location while the communication is ongoing and such a read would retrieve the data written by the application.

Use \texttt{gaspi\_notify} to synchronise the communication.

8.2.2 gaspi\_read

The simplest form of a read operation is \texttt{gaspi\_read} which is a single communication call to read data from a remote location. It is an \textit{asynchronous non-local time-based blocking} procedure.
8.2 Basic communication calls

GASPI_READ (segment_id_local, offset_local, rank, segment_id_remote, offset_remote, size, queue, timeout)

Parameter:

(in) segment_id_local: the local segment ID to write to
(in) offset_local: the local offset in bytes to write to
(in) rank: the remote rank to read from
(in) segment_id_remote: the remote segment to read from
(in) offset_remote: the remote offset to read from
(in) size: the size of the data to read
(in) queue: the queue to use
(in) timeout: the timeout

gaspi_return_t
gaspi_read (gaspi_segment_id_t segment_id_local, gaspi_offset_t offset_local, gaspi_rank_t rank, gaspi_segment_id_t segment_id_remote, gaspi_offset_t offset_remote, gaspi_size_t size, gaspi_queue_id_t queue, gaspi_timeout_t timeout)

function gaspi_read(segment_id_local, offset_local, &rank, segment_id_remote, offset_remote, size, &queue, timeout_ms) &result(res)
integer(gaspi_segment_id_t), value :: segment_id_local
integer(gaspi_offset_t), value :: offset_local
integer(gaspi_rank_t), value :: rank
integer(gaspi_segment_id_t), value :: segment_id_remote
integer(gaspi_offset_t), value :: offset_remote
integer(gaspi_size_t), value :: size
integer(gaspi_queue_id_t), value :: queue
integer(gaspi_timeout_t), value :: timeout_ms
integer(gaspi_return_t) :: res
end function gaspi_read

Execution phase:
Working

Return values:

GASPI_SUCCESS: operation has returned successfully
GASPI_TIMEOUT: operation has run into a timeout
GASPI_ERROR: operation has finished with an error

gaspi_read posts a communication request which asynchronously transfers a contiguous block of size bytes from a source location of a remote GASPI process to a target location of the local GASPI process. This communication request is posted to the communication queue queue. The target location is specified by the pair segment_id_local, offset_local. The source location is specified by the triple rank, segment_id_remote, offset_remote.

A valid gaspi_read communication request requires that the local and the remote segment are allocated, that there is a connection between the local and the remote GASPI process and that the remote segment has been registered on the local GASPI process. Otherwise, the communication request is invalid and the procedure returns with GASPI_ERROR.

After successful procedure completion, i.e. return value GASPI_SUCCESS, the communication request has been posted to the underlying network infrastructure. One new entry is inserted into the given queue.

gaspi_read calls may be posted from every thread of the GASPI process.

If the procedure returns with GASPI_TIMEOUT, the communication request could not be posted to the hardware during the given timeout. This can happen, if another thread is in a gaspi_wait for the same queue. A subsequent call of gaspi_read has to be invoked in order to complete the read call.

A communication request posted to a given queue can be considered as completed, if the the correspondent gaspi_wait returns with GASPI_SUCCESS. For completed gaspi_read requests, the data is guaranteed to be locally available.

If the queue to which the communication request is posted is full, i.e. that the number of posted communication requests has already reached the queue size of a given queue, the communication request fails and the procedure returns with return value GASPI_ERROR. If a saturated queue is detected, there are the following two options: Either one invokes a gaspi_wait on the given queue in order to wait for all the posted requests to be finished. Or one tries to use another queue.
8.2 Basic communication calls

User advice: Return value GASPI_SUCCESS does not mean, that the data transfer has started or that the data has been received at the local side. It is allowed to write data to the local target location while the communication is ongoing. However, the content of the memory would be some undefined interleaving of the data transferred from remote side and the data written locally.
Also, it is allowed to read from the local target location while the communication is ongoing. Such a read would retrieve some undefined interleaving of the data that was present when the call was issued and the data that was transferred from the remote side.

8.2.3 gaspi_wait

The gaspi_wait procedure is a time-based blocking local procedure which waits until all one-sided communication requests posted to a given queue are processed by the network infrastructure. It is an asynchronous non-local time-based blocking procedure.

GASPI_WAIT ( queue , timeout )

Parameter:
(in) queue: the queue ID to wait for
(in) timeout: the timeout

gaspi_wait ( gaspi_queue_id_t queue , gaspi_timeout_t timeout )

function gaspi_wait(queue,timeout_ms) &
&   result( res ) bind(C, name="gaspi_wait")
   integer(gaspi_queue_id_t), value :: queue
   integer(gaspi_timeout_t), value :: timeout_ms
   integer(gaspi_return_t) :: res
end function gaspi_wait

Execution phase:

Working

Return values:
GASPI_SUCCESS: operation has returned successfully
GASPI_TIMEOUT: operation has run into a timeout
GASPI_ERROR: operation has finished with an error
After successful procedure completion, i.e. return value GASPI_SUCCESS, the hitherto posted communication requests have been processed by the network infrastructure and the queue is cleaned up. After that, any communication request which has been posted to the given queue can be considered as completed on the local and remote side. For completed requests the transmitted data is available for the application.

gaspi_wait procedure calls may be posted from every thread of the local Gaspi process. However, the wait operation is a thread exclusive operation and therefore needs privileged access to the queue which means that if a write/read is done while a wait is in operation, the write/read operation blocks to ensure correctness. Enforcing this provides correctness and safety to the user while being easier for the implementor and still allows for a high performance implementation. As a consequence, successive gaspi_wait calls invoked for the same queue by different threads are processed in some sequence one after another.

If the procedure returns with GASPI_TIMEOUT, the wait request could not be completed during the given timeout. This can happen, if there is another thread in a gaspi_wait for the same queue. A subsequent call of gaspi_wait has to be invoked in order to complete the call.

If the procedure returns with GASPI_ERROR, the wait request aborted abnormally.

In both cases, GASPI_TIMEOUT and GASPI_ERROR, the Gaspi state vector should be checked in order to eliminate the possibility of a failure. If a failure is detected, all of the communication requests which have been posted to the given queue since the last gaspi_wait are in an undefined state. Here, undefined state means that the local Gaspi process does not know which requests have been processed and which requests are still outstanding. A call to gaspi_queue_purge has to be invoked in order to reset the queue.

User advice: Return value GASPI_SUCCESS means, that the data of all posted write requests has been transferred to the remote side. It does not mean, that the data has arrived at the remote side. However, write accesses to the local source location will not affect the data that is placed in the remote target location.

User advice: Return value GASPI_SUCCESS means, that the data of all posted read requests have arrived at the local side.

8.2.4 Examples

Listing 11 shows a matrix transpose of a distributed square matrix implemented with the function gaspi_write.

Listing 11: Gaspi all-to-all communication (matrix transpose) implemented with gaspi_write

```c
#include <stdlib.h>
#include <GASPI.h>
#include <success_or_die.h>
```
8.2 Basic communication calls

```c
#include <wait_if_queue_full.h>

extern void dump (int *arr, int nProc);

int
main (int argc, char *argv[])
{
    ASSERT (gaspi_proc_init (GASPI_BLOCK));
    gaspi_rank_t iProc;
    gaspi_rank_t nProc;
    ASSERT (gaspi_proc_rank (&iProc));
    ASSERT (gaspi_proc_num (&nProc));
    gaspi_notification_id_t notification_max;
    ASSERT (gaspi_notification_num(&notification_max));

    if (notification_max < (gaspi_notification_id_t)nProc)
    {
        exit (EXIT_FAILURE);
    }
    ASSERT (gaspi_group_commit (GASPI_GROUP_ALL, GASPI_BLOCK));
    const gaspi_segment_id_t segment_id_src = 0;
    const gaspi_segment_id_t segment_id_dst = 1;
    const gaspi_size_t segment_size = nProc * sizeof(int);
    ASSERT (gaspi_segment_create ( segment_id_src, segment_size,
        GASPI_GROUP_ALL, GASPI_BLOCK,
        GASPI_ALLOC_DEFAULT ));
    ASSERT (gaspi_segment_create ( segment_id_dst, segment_size,
        GASPI_GROUP_ALL, GASPI_BLOCK,
        GASPI_ALLOC_DEFAULT ));

    int *src = NULL;
    int *dst = NULL;
    ASSERT (gaspi_segment_ptr (segment_id_src, &src));
    ASSERT (gaspi_segment_ptr (segment_id_dst, &dst));

    for (gaspi_rank_t rank = 0; rank < nProc; ++rank)
    {
```
8.2 Basic communication calls

```c
{
    src[rank] = iProc * nProc + rank;

    const gaspi_offset_t offset_src = rank * sizeof (int);
    const gaspi_offset_t offset_dst = iProc * sizeof (int);
    const gaspi_notification_id_t notify_ID = rank;

    wait_if_queue_full (queue_id, 2);

    const gaspi_notification_t notify_val = 1;

    ASSERT (gaspi_write_notify (segment_id_src, offset_src,
                                rank, segment_id_dst, offset_dst,
                                sizeof (int), notify_ID, notify_val,
                                queue_id, GASPI_BLOCK));
}

gaspi_notification_id_t notify_cnt = nProc;
gaspi_notification_id_t first_notify_id;

while (notify_cnt > 0)
{
    ASSERT (gaspi_notify_waitsome (segment_id_dst, 0, nProc,
                                    &first_notify_id, GASPI_BLOCK));

    gaspi_notification_id_t notify_val = 0;
    ASSERT (gaspi_notify_reset (segment_id_dst, first_notify_id,
                                &notify_val));

    if (notify_val != 0)
    {
        --notify_cnt;
    }
}
dump (dst, nProc);

ASSERT (gaspi_wait (queue_id, GASPI_BLOCK));

ASSERT (gaspi_barrier (GASPI_GROUP_ALL, GASPI_BLOCK));

ASSERT (gaspi_proc_term (GASPI_BLOCK));

return EXIT_SUCCESS;
}
```
Listing 12 shows a matrix transpose of a distributed square matrix implemented with the function \texttt{gaspi\_read}. Please note the differences between the transpose implemented with write and the transpose implemented with read: The implementation using write can initialize the matrix on-the-fly, right before the data is transferred, while the implementation using read has to synchronise all processes after the local initialization in order to be sure to read valid data. On the other hand, in the implementation using write one has to synchronise after the local wait whereas in the implementation using read one can directly use the data after the local wait returns.

Listing 12: \texttt{GASPI} all-to-all communication (matrix transpose) implemented with \texttt{gaspi\_read}

```c
#include <stdlib.h>
#include <GASPI.h>
#include <success_or_die.h>
#include <wait_if_queue_full.h>

extern void dump (int *arr, int nProc);

int main (int argc, char *argv[])
{
    ASSERT (gaspi_proc_init (GASPI_BLOCK));

gaspi_rank_t iProc;
gaspi_rank_t nProc;

    ASSERT (gaspi_proc_rank (&iProc));
    ASSERT (gaspi_proc_num (&nProc));

    ASSERT (gaspi_group_commit (GASPI_GROUP_ALL, GASPI_BLOCK));

    const gaspi_segment_id_t segment_id_src = 0;
    const gaspi_segment_id_t segment_id_dst = 1;

    const gaspi_size_t segment_size = nProc * sizeof(int);

    ASSERT (gaspi_segment_create ( segment_id_src, segment_size ,
                                    GASPI_GROUP_ALL, GASPI_BLOCK ,
                                    GASPI_ALLOC_DEFAULT )
             );

    ASSERT (gaspi_segment_create ( segment_id_dst, segment_size ,
                                    GASPI_GROUP_ALL, GASPI_BLOCK ,
                                    GASPI_ALLOC_DEFAULT )
             );

    int *src = NULL;
    int *dst = NULL;
```
8.3 Weak synchronisation primitives

8.3.1 Introduction

The one-sided communication procedures have the characteristics that the entire communication is managed by the local process only. The remote process is not involved. This has the advantage that there is no inherent synchronisa-
8.3 Weak synchronisation primitives

Synchronization between the local and the remote process in every communication request. However, at some point, the remote process needs the information as to whether the data which has been sent to that process has arrived and is valid.

To this end GASPI provides so-called weak synchronisation primitives which allows the application to inform the remote side that the data has been transferred by updating a notification on the remote side. These notifications must be submitted to the same queue to which the data payload has been attached. Otherwise, causality is not guaranteed.

As counterpart, there are routines which wait for an update of a single or even an entire set of notifications. There is a thread safe atomic function to reset the local notification with a given ID which returns the value of the notification before it is reset.

These notification procedures are also one-sided and involve only the local process.

8.3.2 gaspi_notify

gaspi_notify is an asynchronous non-local time-based blocking procedure.

\[
\text{GASPI NOTIFY} ( \text{segment}\_id, \text{rank}, \text{notification}\_id, \text{notification}\_value, \text{queue}, \text{timeout})
\]

**Parameter:**

(in) `segment_id`: the remote segment bound to the notification

(in) `rank`: the remote rank to notify

(in) `notification_id`: the remote notification ID

(in) `notification_value`: the notification value (> 0) to write

(in) `queue`: the queue to use

(in) `timeout`: the timeout

\[
\text{gaspi}\_\text{return} \_t \\
\text{gaspi}\_\text{notify} ( \text{gaspi}\_\text{segment}\_id\_t \text{segment}\_id, \text{gaspi}\_\text{rank}\_t \text{rank}, \text{gaspi}\_\text{notification}\_id\_t \text{notification}\_id, \text{gaspi}\_\text{notification}\_t \text{notification}\_value, \text{gaspi}\_\text{queue}\_id\_t \text{queue}, \text{gaspi}\_\text{timeout}\_t \text{timeout})
\]
8.3 Weak synchronisation primitives

function gaspi_notify(segment_id_remote, rank, notification_id, &
   notification_value, queue, timeout_ms) &
   result(res) bind(C, name="gaspi_notify")
   integer(gaspi_segment_id_t), value :: segment_id_remote
   integer(gaspi_rank_t), value :: rank
   integer(gaspi_notification_id_t), value :: notification_id
   integer(gaspi_notification_t), value :: notification_value
   integer(gaspi_queue_id_t), value :: queue
   integer(gaspi_timeout_t), value :: timeout_ms
   integer(gaspi_return_t) :: res
end function gaspi_notify

Execution phase:

Working

Return values:
GASPI_SUCCESS: operation has returned successfully
GASPI_TIMEOUT: operation has run into a timeout
GASPI_ERROR: operation has finished with an error

gaspi_notify posts a notification request which asynchronously transfers the
notification notification_value of the local GASPI process to an internal notifica-
tion buffer of a remote GASPI process. This notification request is posted to
the communication queue queue. The remote notification buffer is specified by
the pair rank notification_id.

A valid gaspi_notify communication request requires that there is a connection
between the local and the remote GASPI process. Otherwise, the communication
request is invalid and the procedure returns with GASPI_ERROR.

After successful procedure completion, i.e. return value GASPI_SUCCESS, the
notification request has been posted to the underlying network infrastructure.

A gaspi_notify call which is posted subsequent to an arbitrary number of
gaspi_write requests and which is posted to the same queue and the same
destination rank is guaranteed to be non-overtaking. Non-overtaking means
that the order of communication requests is preserved on the remote side. In
particular, one can assume, that if the data from the gaspi_notify request
has arrived on the remote process, also the data from the earlier posted write
request(s) to the same process have arrived on the remote side.

gaspi_notify calls may be posted from every thread of the GASPI process.

If the procedure returns with GASPI_TIMEOUT, the notification request could
not be posted to the hardware during the given timeout. This can happen if
another thread is in a gaspi_wait for the same queue. A subsequent call of
gaspi_notify has to be invoked in order to complete the call.

A notification request posted to a given queue can be considered as completed,
if the the correspondent gaspi_wait returns with GASPI_SUCCESS.
8.3 Weak synchronisation primitives

If the queue to which the communication request is posted is full, i.e. that the number of posted communication requests has already reached the queue size of a given queue, the communication request fails.

User advice: Return value GASPI_SUCCESS does not mean, that the notification has been transferred or that the notification has arrived at the remote side.

8.3.3 gaspi_notify_waitsome

For the procedures with notification, gaspi_notify and the extendend function gaspi_write_notify, gaspi_notify_waitsome is the correspondent wait procedure for the receiver (notified) side. gaspi_notify_waitsome is a synchronous, non-local time-based blocking procedure.

**GASPI_NOTIFY_WAITSOME (**

- segment_id:
- notification_begin:
- notification_num:
- first_id:
- timeout

**Parameter:**

(in) segment_id: the segment bound to the notification

(in) notification_begin: the local notification ID for the first notification to wait for

(in) notification_num: the number of notification ID’s to wait for

(out) first_id: the id of the first notification that arrived

(in) timeout: the timeout

**Function gaspi_notify_waitsome(segment_id_local,&**

& notification_begin,num,first_id,timeout_ms) &

& result ( res ) bind(C, nam="gaspi_notify_waitsome")

integer(gaspi_segment_id_t), value :: segment_id_local
integer(gaspi_notification_id_t), value :: notification_begin
integer(gaspi_number_t), value :: num
integer(gaspi_notification_id_t) :: first_id
integer(gaspi_timeout_t), value :: timeout_ms
integer(gaspi_return_t) :: res
end function gaspi_notify_waitsome
8.3 Weak synchronisation primitives

Execution phase:
Working

Return values:
GASPI_SUCCESS: operation has returned successfully
GASPI_TIMEOUT: operation has run into a timeout
GASPI_ERROR: operation has finished with an error

gaspi_notify_waitsome waits that at least one of a number of consecutive notifications residing in the local internal buffer has a value that is not zero.
The notification buffer is specified by the pair notification_begin, notification_num. It contains notification_num many consecutive notifications beginning at the notification with ID notification_begin.
If notification_num == 0 then gaspi_notify_waitsome returns immediately with GASPI_SUCCESS.
After successful procedure completion, i.e. return value GASPI_SUCCESS, the value of at least one of the notifications in the notification buffer has changed to a value that is not zero. All threads that are waiting for the notifications are notified.
If the procedure returns with GASPI_TIMEOUT, no notification has changed during the given period of time.
In case of an error, i.e. GASPI_ERROR, the values of the notifications are undefined.

User advice: One scenario for the usage of gaspi_notify_waitsome inspecting only one notification is the following: The remote side uses a gaspi_write call followed by a subsequent call of gaspi_notify posted to the same queue and the same destination rank. GASPI guarantees, that if the notification has arrived on the remote process, the previously posted request carrying the work load has arrived as well.

User advice: If in a multi-threaded application more than one thread calls gaspi_notify_waitsome for the range of notifications, then all waiting threads are notified about the change of at least one of the notifications. By inspecting the actual values of each of the notifications with gaspi_notify_reset, only one thread per changed notification receives a value different from zero.

User advice: In a multi-threaded application the code in listing 13 selects one thread to act on the change of a single notification. The code waits in a blocking manner and thus cannot be used in failure tolerant applications.

Listing 13: Blocking waitsome in a multi-threaded application
8.3 Weak synchronisation primitives

```c
#include <GASPI.h>
#include <success_or_die.h>

extern void process ( const gaspi_notification_id_t id,
                      const gaspi_notification_t val );

void blocking_waitsome ( const gaspi_notification_id_t id_begin,
                         const gaspi_notification_id_t id_end,
                         const gaspi_segment_id_t seg_id )
{
    gaspi_notification_id_t first_id;
    ASSERT ( gaspi_notify_waitsome ( seg_id,
                                    id_begin,
                                    id_end - id_begin,
                                    &first_id,
                                    GASPI_BLOCK ) );
    gaspi_notification_t val = 0;
    // atomic reset
    ASSERT ( gaspi_notify_reset ( seg_id, first_id, &val ) );
    // other threads are notified too!
    process ( first_id, val );
}
```

8.3.4 gaspi_notify_reset

For the `gaspi_notify_waitsome` procedure, there is a notification initialization procedure which resets the given notification to zero. It is a synchronous local blocking procedure.

```
GASPI_NOTIFY_RESET ( segment_id
                     , notification_id
                     , old_notification_val )
```

Parameter:

*(in) segment_id:* the segment bound to the notification

*(in) notification_id:* the local notification ID to reset

*(out) old_notification_val:* notification value before reset
8.4 Extended communication calls

```fortran
function gaspi_notify_reset(segment_id_local, &
&      notification_id, old_notification_val) &
&      result(res) bind(C, name="gaspi_notify_reset")
  integer(gaspi_segment_id_t), value :: segment_id_local
  integer(gaspi_notification_id_t), value :: notification_id
  integer(gaspi_notification_t) :: old_notification_val
  integer(gaspi_return_t) :: res
  end function gaspi_notify_reset
```

**Execution phase:**

**Working**

**Return values:**

GASPI_SUCCESS: operation has returned successfully

GASPI_ERROR: operation has finished with an error

**gaspi_notify_reset** resets the notification with ID **notification_id** to zero. The function **gaspi_notify_reset** is an atomic operation: Threads can use **gaspi_notify_reset** to safely extract the value of a specific notification.

The notification buffer on the local side is specified by the notification ID **notification_id**.

After successful procedure completion, i.e. return value GASPI_SUCCESS, the value of the notification buffer was set to zero and **old_notification_val** contains the content of the notification buffer before it was set to zero. To read the old value and to set the value to zero is a single atomic operation.

**gaspi_notify_reset** calls may be posted from every thread of the Gaspi process.

In case of error, i.e. return value GASPI_ERROR, the value of **old_notification_val** is undefined.

### 8.4.1 gaspi_write_notify

The **gaspi_write_notify** variant extends the simple **gaspi_write** with a notification on the remote side. This applies to communication patterns that require
8.4 Extended communication calls

tighter synchronisation on data movement. The remote receiver of the data is
notified when the write is finished and can verify this through the respective
wait procedure. It is an asynchronous non-local time-based blocking procedure.

GASPI_WRITE_NOTIFY ( segment_id_local
    , offset_local
    , rank
    , segment_id_remote
    , offset_remote
    , size
    , notification_id
    , notification_value
    , queue
    , timeout )

Parameter:

(in) segment_id_local: the local segment ID to read from
(in) offset_local: the local offset in bytes to read from
(in) rank: the remote rank to write to
(in) segment_id_remote: the remote segment to write to
(in) offset_remote: the remote offset to write to
(in) size: the size of the data to write
(in) notification_id: the remote notification ID
(in) notification_value: the value of the notification to write
(in) queue: the queue to use
(in) timeout: the timeout

gaspi_return_t
gaspi_write_notify ( gaspi_segment_id_t segment_id_local
    , gaspi_offset_t offset_local
    , gaspi_rank_t rank
    , gaspi_segment_id_t segment_id_remote
    , gaspi_offset_t offset_remote
    , gaspi_size_t size
    , gaspi_notification_id_t notification_id
    , gaspi_notification_t notification_value
    , gaspi_queue_id_t queue
    , gaspi_timeout_t timeout )
8.4 Extended communication calls

function gaspi_write_notify(segment_id_local, offset_local, & rank, segment_id_remote, offset_remote, size, & notification_id, notification_value, queue, & timeout_ms) &
& result( res ) bind(C, name="gaspi_write_notify")
integer(gaspi_segment_id_t), value :: segment_id_local
integer(gaspi_offset_t), value :: offset_local
integer(gaspi_rank_t), value :: rank
integer(gaspi_segment_id_t), value :: segment_id_remote
integer(gaspi_offset_t), value :: offset_remote
integer(gaspi_size_t), value :: size
integer(gaspi_notification_id_t), value :: notification_id
integer(gaspi_notification_t), value :: notification_value
integer(gaspi_queue_id_t), value :: queue
integer(gaspi_timeout_t), value :: timeout_ms
integer(gaspi_return_t) :: res
end function gaspi_write_notify

Execution phase:
Working

Return values:
GASPI_SUCCESS: operation has returned successfully
GASPI_TIMEOUT: operation has run into a timeout
GASPI_ERROR: operation has finished with an error

Implementor advice: The procedure is semantically equivalent to a call to gaspi_write and a subsequent call of gaspi_notify. However, it should be implemented more efficiently, if supported by the network infrastructure.

8.4.2 gaspi_write_list

The gaspi_write_list variant allows strided communication where a list of different data locations are processed at once. Semantically, it is equivalent to a sequence of calls to gaspi_write but it should (if possible) be more efficient. It is an asynchronous non-local time-based blocking procedure.

GASPI_WRITE_LIST ( num
, segment_id_local[num]
, offset_local[num]
, rank
, segment_id_remote[num]
, offset_remote[num]
, size[num]
, queue
, timeout )
8.4 Extended communication calls

Parameter:

(in) num: the number of elements to write
(in) segment_id_local[num]: list of local segment ID’s to read from
(in) offset_local[num]: list of local offsets in bytes to read from
(in) rank: the remote rank to write to
(in) segment_id_remote[num]: list of remote segments to write to
(in) offset_remote[num]: list of remote offsets to write to
(in) size[num]: list of sizes of the data to write
(in) queue: the queue to use
(in) timeout: the timeout

```fortran
function gaspi_write_list(num, segment_id_local, offset_local, & rank, segment_id_remote, offset_remote, size, queue, & timeout_ms) &
  result(res) bind(C, name="gaspi_write_list")
  integer(gaspi_number_t), value :: num
  type(c_ptr), value :: segment_id_local
  type(c_ptr), value :: offset_local
  integer(gaspi_rank_t), value :: rank
  type(c_ptr), value :: segment_id_remote
  type(c_ptr), value :: offset_remote
  type(c_ptr), value :: size
  integer(gaspi_queue_id_t), value :: queue
  integer(gaspi_timeout_t), value :: timeout_ms
  integer(gaspi_return_t) :: res
end function gaspi_write_list
```

Execution phase:

Working

Return values:

GASPI_SUCCESS: operation has returned successfully
8.4 Extended communication calls

GASPI_TIMEOUT: operation has run into a timeout
GASPI_ERROR: operation has finished with an error

Implementor advice: The procedure is semantically equivalent to num subsequent calls of gaspi_write with the given local and remote location specification, provided that the destination rank and the used queue are invariant. However, it should be implemented more efficiently, if supported by the network infrastructure.

8.4.3 gaspi_write_list_notify

The gaspi_write_list_notify operation performs strided communication as gaspi_write_list but also includes a notification that the remote receiver can use to ensure that the communication step is completed. It is an asynchronous non-local time-based blocking procedure.

\[
\text{GASPI\_WRITE\_LIST\_NOTIFY} \\
\geq \begin{array}{l}
\text{num, segment_id_local[num], offset_local[num], rank, segment_id_remote[num], offset_remote[num], size[num], notification_id, notification_value, queue, timeout}
\end{array}
\]

Parameter:

\(\text{(in) num: the number of elements to write}\)
\(\text{(in) segment_id_local[num]: list of local segment ID’s to read from}\)
\(\text{(in) offset_local[num]: list of local offsets in bytes to read from}\)
\(\text{(in) rank: the remote rank to be write to}\)
\(\text{(in) segment_id_remote[num]: list of remote segments to write to}\)
\(\text{(in) offset_remote[num]: list of remote offsets to write to}\)
\(\text{(in) size[num]: list of sizes of the data to write}\)
\(\text{(in) notification_id: the remote notification ID}\)
\(\text{(in) notification_value: the value of the notification to write}\)
\(\text{(in) queue: the queue to use}\)
\(\text{(in) timeout: the timeout}\)
8.4 Extended communication calls

```fortran
function gaspi_write_list_notify(num, segment_id_local, &
  offset_local, rank, segment_id_remote, &
  offset_remote, size, segment_id_notification, &
  notification_id, notification_value, queue, timeout_ms) &
  result(res) bind(C, name="gaspi_write_list_notify")
  integer(gaspi_number_t), value :: num
  type(c_ptr), value :: segment_id_local
  type(c_ptr), value :: offset_local
  integer(gaspi_rank_t), value :: rank
  type(c_ptr), value :: segment_id_remote
  type(c_ptr), value :: offset_remote
  type(c_ptr), value :: size
  integer(gaspi_segment_id_t), value :: segment_id_notification
  integer(gaspi_notification_id_t), value :: notification_id
  integer(gaspi_notification_t), value :: notification_value
  integer(gaspi_queue_id_t), value :: queue
  integer(gaspi_timeout_t), value :: timeout_ms
  integer(gaspi_return_t) :: res
end function gaspi_write_list_notify
```

**Execution phase:**

Working

**Return values:**

GASPI\_SUCCESS: operation has returned successfully

GASPI\_TIMEOUT: operation has run into a timeout

GASPI\_ERROR: operation has finished with an error

**Implementor advice:** The procedure is semantically equivalent to a call to gaspi\_write\_list and a subsequent call of gaspi\_notify. However, it should be implemented more efficiently, if supported by the network infrastructure.
8.4 Extended communication calls

8.4.4 gaspi_read_list

The `gaspi_read_list` variant allows strided communication where a list of different data locations are processed at once. Semantically, it is equivalent to a sequence of calls to `gaspi_read` but it should (if possible) be more efficient. It is an asynchronous non-local time-based blocking procedure.

```
GASPI_READ_LIST ( num
    , segment_id_local[num]
    , offset_local[num]
    , rank
    , segment_id_remote[num]
    , offset_remote[num]
    , size[num]
    , queue
    , timeout )
```

**Parameter:**

*(in) num:* the number of elements to read
*(in) segment_id_local[num]:* list of local segment ID’s to write to
*(in) offset_local[num]:* list of local offsets in bytes to write to
*(in) rank:* the remote rank to read from
*(in) segment_id_remote[num]:* list of remote segments to read from
*(in) offset_remote[num]:* list of remote offsets to read from
*(in) size[num]:* list of sizes of the data to read
*(in) queue:* the queue to use
*(in) timeout:* the timeout

```
gaspi_return_t gaspi_read_list ( gaspi_number_t num
    , gaspi_segment_id_t const *segment_id_local
    , gaspi_offset_t const *offset_local
    , gaspi_rank_t rank
    , gaspi_segment_id_t const *segment_id_remote
    , gaspi_offset_t const *offset_remote
    , gaspi_size_t const *size
    , gaspi_queue_id_t queue
    , gaspi_timeout_t timeout )
```
8.5 Communication utilities

8.5.1 gaspi_queue_size

The gaspi_queue_size procedure is a *synchronous local blocking* procedure which determines the number of open communication requests posted to a given queue.

\[
\text{GASPI\_QUEUE\_SIZE} \ (queue, \ queue\_size )
\]

*Parameters:*

*(in) queue:* the queue to probe

*(out) queue_size:* the number of open requests posted to the queue

\[
gaspi\_return\_t
gaspi_queue_size \ (gaspi\_queue\_id\_t \ queue,
\quad gaspi\_number\_t \ const \ *queue\_size )
\]
8.5 Communication utilities

function gaspi_queue_size(queue, queue_size) &
&   result(res) bind(C, name="gaspi_queue_size")
   integer(gaspi_queue_id_t), value :: queue
   integer(gaspi_number_t) :: queue_size
   integer(gaspi_return_t) :: res
end function gaspi_queue_size

Execution phase:
Working

Return values:
GASPI_SUCCESS: operation has returned successfully
GASPI_ERROR: operation has finished with an error

After successful procedure completion, i.e. return value GASPI_SUCCESS, the parameter queue_size contains the number of open requests posted to the queue queue. In a threaded program this result is uncertain, since another thread may have posted an additional request in the meantime or issued a wait call.

The queue size is set to zero by a successful call to gaspi_wait.

In case of error, the return value is GASPI_ERROR. The parameter queue_size has an undefined value.

8.5.2 gaspi_queue_purge

The gaspi_queue_purge procedure is a synchronous local time-based blocking procedure which purges a given queue.

GASPI_QUEUE_PURGE (queue, timeout)

Parameter:
(in) queue: the queue to purge
(in) timeout: the timeout

gaspi_return_t

gaspi_queue_purge (gaspi_queue_id_t queue,
      gaspi_timeout_t timeout)

function gaspi_queue_purge(queue, timeout) &
&   result(res) bind(C, name="gaspi_queue_purge")
   integer(gaspi_queue_id_t), value :: queue
   integer(gaspi_timeout_t), value :: timeout
   integer(gaspi_return_t) :: res
end function gaspi_queue_purge
Execution phase:
Working

Return values:
GASPI_SUCCESS: operation has returned successfully
GASPI_TIMEOUT: operation has run into a timeout
GASPI_ERROR: operation has finished with an error

This procedure should only be invoked in the situation in which a node failure is detected by inspecting the global health state with gaspi_state_vec_get.

After successful procedure completion, i.e. return value GASPI_SUCCESS, the communication queue is purged. All communication requests posted to the queue are eliminated from the queue. The local GASPI process has no information about the completion of communication requests posted to the given queue since the last invocation of gaspi_wait.

If the procedure returns with GASPI_TIMEOUT, the purge request could not be completed during the given timeout. This might happen if there is another thread in a gaspi_wait for the same queue. A subsequent call of gaspi_queue_purge has to be invoked in order to complete the call.

If the procedure returns with GASPI_ERROR, the purge request aborted abnormally.

9 Passive communication

9.1 Introduction and overview

Passive communication has a two-sided semantic, where there is a matching receiver to a send request. Passive communication aims at communication patterns where the sender is unknown (i.e. it can be any process from the receiver perspective) but there is potentially the need for synchronisation between processes. Typical example uses cases are:

- Distributed update where many processes contribute to the data of one process.
- Pass arguments and results.
- Global error handling.

The implementation should try to enforce fairness in communication that is, no sender should see its communication request delayed indefinitely.

The passive keyword means that the communication calls should avoid busy-waiting and consume no CPU cycles, freeing the system for computation.
Both the send and the matching receive are time-based blocking. A valid passive communication request requires that the local and the remote segment are allocated and that there is a connection between the local and the remote GASPI process. Otherwise, the communication request is invalid and the procedure returns with GASPI_ERROR.

### 9.2 Passive communication calls

#### 9.2.1 gaspi_passive_send

The blocking gaspi_passive_send is the routine called by the sender side to engage in passive communication. It is a synchronous non-local time-based blocking procedure.

```plaintext
GASPI_PASSIVE_SEND ( segment_id_local
, offset_local
, rank
, size
, timeout )
```

**Parameter:**

- *(in) segment_id_local:* the local segment ID from which the data is sent
- *(in) offset_local:* the local offset from which the data is sent
- *(in) rank:* the remote rank to which the data is sent
- *(in) size:* the size of the data to be sent
- *(in) timeout:* the timeout

```plaintext
function gaspi_passive_send ( segment_id_local, offset_local, rank, size, timeout_ms )
& result( res ) bind(C, name="gaspi_passive_send")
integer(gaspi_segment_id_t), value :: segment_id_local
integer(gaspi_offset_t), value :: offset_local
integer(gaspi_rank_t), value :: rank
integer(gaspi_size_t), value :: size
integer(gaspi_timeout_t), value :: timeout_ms
integer(gaspi_return_t) :: res
end function gaspi_passive_send
```

**Execution phase:**
9.2 Passive communication calls

Working

Return values:
GASPI_SUCCESS: operation has returned successfully
GASPI_TIMEOUT: operation has run into a timeout
GASPI_ERROR: operation has finished with an error

gaspi_passive_send posts a passive communication request which transfers a contiguous block of size bytes from a source location of the local GASPI process to the remote GASPI process with the indicated rank rank. On the remote side, a corresponding gaspi_passive_receive has to be posted. The source location is specified by the pair segment_id_local, offset_local.

There is a size limit for the data sent with gaspi_passive_send. The maximum size is returned by the function gaspi_passive_transfer_size_max.

A valid gaspi_passive_send communication request requires that the local and the remote segment are allocated and that there is a connection between the local and the remote GASPI process. Otherwise, the communication request is invalid and the procedure returns with GASPI_ERROR.

After successful procedure completion, i.e. return value GASPI_SUCCESS, the passive communication request has been posted to the underlying network infrastructure and was completed.

gaspi_passive_send calls may be posted from every thread of the GASPI process.

If the procedure returns with GASPI_TIMEOUT, the communication request could not be posted to the hardware during the given timeout.

If the passive communication queue is full at the time when a new passive communication request is posted, i.e. the number of posted communication requests has already reached the queue size, the communication request fails and the procedure returns with return value GASPI_ERROR.

User advice: Since the passive receive will try to match every corresponding send, the buffer sizes for send/recv need to match for all ranks for the passive communication within one passive send/recv communication step.

User advice:[see also the advice in 8.2.1 on page 56] It is allowed to write data to the source location while the communication is ongoing. However, the result on the remote side would be some undefined interleaving of the data that was present when the call was issued and the data that was written later.
It is also allowed to read from the source location while the communication is ongoing and such a read would retrieve the data written by the application.
User advice: If the parameter `build Infrastructure` is not set, a connection has to be established between the processes before the `gaspi_passive_send` can be used. This is accomplished calling the procedure `gaspi_connect`.

9.2.2 `gaspi_passive_receive`

The synchronous non-local time-based blocking `gaspi_passive_receive` is one of the routines called by the receiver side to engage in passive communication.

```c
GASPI_PASSIVE_RECEIVE ( segment_id_local
    , offset_local
    , rank
    , size
    , timeout )
```

Parameter:

- **(in) segment_id_local**: the local segment ID where to write the data
- **(in) offset_local**: the local offset where to write the data
- **(out) rank**: the remote rank from which the data is transferred
- **(in) size**: the size of the data to be received
- **(in) timeout**: the timeout

```c
function gaspi_passive_receive(segment_id_local,offset_local, &
    & rem_rank,size,timeout_ms) &
    & result( res ) bind(C, name="gaspi_passive_receive")
integer(gaspi_segment_id_t), value :: segment_id_local
integer(gaspi_offset_t), value :: offset_local
integer(gaspi_rank_t) :: rem_rank
integer(gaspi_size_t), value :: size
integer(gaspi_timeout_t), value :: timeout_ms
integer(gaspi_return_t) :: res
end function gaspi_passive_receive
```

Execution phase:

Working

Return values:
9.2 Passive communication calls

**GASPI_SUCCESS**: operation has returned successfully

**GASPI_TIMEOUT**: operation has run into a timeout

**GASPI_ERROR**: operation has finished with an error

gaspi_passive_receive receives a contiguous block of data into a target location from some unspecified remote Gaspi process. The target location is specified by the pair `segment_id_local, offset_local`.

There is no need for the `gaspi_passive_receive` procedure to be active before a corresponding `gaspi_passive_send` procedure is invoked. However, as long as there is no matching receive, the `gaspi_passive_send` cannot achieve any progress and thus cannot return `GASPI_SUCCESS`.

The target location needs to have enough space to hold the maximum passive transfer size that could be sent by any other process. Otherwise, the received data might overwrite memory regions outside of the allocated memory and the application will be in an undefined state.

A valid `gaspi_passive_receive` communication request requires that the local destination segment is allocated and that there is a connection between the local and the remote Gaspi process from which a data transfer originates. Otherwise, the communication request is invalid and the procedure returns with **GASPI_ERROR**.

After successful procedure completion, i.e. return value **GASPI_SUCCESS**, the data has been received and is available at the target location. Further `rank` contains the rank of the sending process and associated to the communication request.

Successive `gaspi_passive_receive` calls posted by two different threads using two different target locations are allowed. However, the first incoming data is received either by the first thread or the by the second. That means that the `gaspi_passive_receive` should be posted only from a single thread of a Gaspi process.

If the procedure returns with **GASPI_TIMEOUT**, there was no pending communication request in the queue. The output parameter `rank` has no defined value.

*User advice*: It is allowed to write data to the local target location while the passive communication is ongoing. However, the content of the memory would be some undefined interleaving of the data transferred from remote side and the data written locally.

Also, it is allowed to read from the local target location while the passive communication is ongoing. Such a read would retrieve some undefined interleaving of the data that was present when the call was issued and the data that was transferred from the remote side.

*Implementor advice*: A quality implementation enforces fairness in communication that is, no sender should see its communication request delayed indefinitely. The passive keyword means the communication calls shall avoid busy-waiting and consume no CPU cycles, freeing the system for computation.
9.3 Passive communication utilities

9.3.1 `gaspi_passive_queue_purge`

The `gaspi_passive_queue_purge` procedure is a *synchronous local time-based blocking* procedure which purges the passive queue.

```
GASPI_PASSIVE_QUEUE_PURGE (timeout)
```

**Parameter:**

*(in) timeout:* the timeout

```
gaspi_return_t
gaspi_passive_queue_purge (gaspi_timeout_t timeout)
```

```
function gaspi_passive_queue_purge(timeout) &
& result(res) bind(C, name="gaspi_passive_queue_purge")
  integer(gaspi_timeout_t), value :: timeout
  integer(gaspi_return_t) :: res
end function gaspi_passive_queue_purge
```

**Execution phase:**

Working

**Return values:**

- **GASPI_SUCCESS**: operation has returned successfully
- **GASPI_TIMEOUT**: operation has run into a timeout
- **GASPI_ERROR**: operation has finished with an error

This procedure should only be invoked in the situation in which a node failure is detected by inspecting the global health state with `gaspi_state_vec_get`.

After successful procedure completion, i.e. return value **GASPI_SUCCESS**, the passive communication queue is purged.

If the procedure returns with **GASPI_TIMEOUT**, the purge request could not be completed during the given timeout. A subsequent call of `gaspi_passive_queue_purge` has to be invoked in order to complete the call.

If the procedure returns with **GASPI_ERROR**, the purge request was not satisfied and returned abnormally.
10 Global atomics

10.1 Introduction and Overview

An atomic operation is an operation which is guaranteed to be executed without fear of interference from other processes during the procedure call. Only one GASPI process at a time has access to the global variable and can modify it. Atomic operations are also guaranteed to be fair. That means no GASPI process should see its atomic operation request delayed indefinitely.

10.2 Atomic operation calls

10.2.1 gaspi_atomic_fetch_add

The gaspi_atomic_fetch_add procedure is a synchronous non-local time-based blocking procedure which atomically adds a given value to a globally accessible value.

\[
\text{gaspi_return_t gaspi_atomic_fetch_add ( gaspi_segment_id_t segment_id, gaspi_offset_t offset, gaspi_rank_t rank, gaspi_atomic_value_t value_add, gaspi_atomic_value_t *value_old, gaspi_timeout_t timeout )}
\]

Parameter:

(in) segment_id: the segment ID where the value is located

(in) offset: the offset where the value is located

(in) rank: the rank where the value is located

(in) value_add: the value which is to be added

(out) value_old: the old value before the operation

(in) timeout: the timeout
function gaspi_atomic_fetch_add(segment_id, offset, rank, &
val_add, val_old, timeout_ms) &
& result(res) bind(C, name="gaspi_atomic_fetch_add")
integer(gaspi_segment_id_t), value :: segment_id
integer(gaspi_offset_t), value :: offset
integer(gaspi_rank_t), value :: rank
integer(gaspi_atomic_value_t), value :: val_add
integer(gaspi_atomic_value_t) :: val_old
integer(gaspi_timeout_t), value :: timeout_ms
integer(gaspi_return_t) :: res
end function gaspi_atomic_fetch_add

Execution phase:

Working

Return values:
GASPI_SUCCESS: operation has returned successfully
GASPI_TIMEOUT: operation has run into a timeout
GASPI_ERROR: operation has finished with an error

gaspi_atomic_fetch_add atomically adds the value of \texttt{value\_add} to the value
on rank \texttt{rank}, segment \texttt{segment\_id\_remote} and offset \texttt{offset\_remote}.

After successful procedure completion, i.e. return value GASPI_SUCCESS, the
parameter \texttt{value\_old} contains the value before the operation has been applied.

If the procedure returns with GASPI_TIMEOUT, the fetch and add request could
not be completed during the given timeout. The parameter \texttt{value\_old} has an
undefined value. A subsequent call of gaspi_atomic_fetch_add needs to be
invoked in order to complete the operation.

If the procedure returns with GASPI_ERROR, the fetch and add request aborted
abnormally. The parameter \texttt{value\_old} as well as the global value (segment_id,
offset, rank) have undefined values.

In both cases, GASPI_TIMEOUT and GASPI_ERROR, the Gaspi state vector should
be checked in order to deal with possible failures.

Implementor advice: The implementation might require some alignment
restrictions that is, the triple(\texttt{segment\_id, offset, rank}) might be re-
quired to respect some alignment restrictions.

User advice: Concurrent accesses to the location represented by the
triple(\texttt{segment\_id, offset, rank}) are possible but consistency must be
handled by the application.

10.2.2 gaspi_atomic_compare_swap

The gaspi_atomic_compare_swap procedure is a synchronous non-local time-
based blocking procedure which atomically compares the value of a global value
against some user given value and in case these are equal the old value is replaced by a new value.

\[
\text{GASPI_ATOMIC_COMPARE_SWAP ( segment_id, offset, rank, comparator, value_new, value_old, timeout )}
\]

**Parameter:**

*(in) segment_id:* the segment ID where the value is located

*(in) offset:* the offset where the value is located

*(in) rank:* the rank where the value is located

*(in) comparator:* the value which is compared to the remote value

*(in) value_new:* the new value to which the remote location is set if the result of the comparison is true

*(out) value_old:* the value before the operation

*(in) timeout:* the timeout

```plaintext
gaspi_return_t gaspi_atomic_compare_swap ( gaspi_segment_id_t segment_id, gaspi_offset_t offset, gaspi_rank_t rank, gaspi_atomic_value_t comparator, gaspi_atomic_value_t value_new, gaspi_atomic_value_t *value_old, gaspi_timeout_t timeout )
```

**Execution phase:**

*Working*
10.2 Atomic operation calls

Return values:
GASPI_SUCCESS: operation has returned successfully
GASPI_TIMEOUT: operation has run into a timeout
GASPI_ERROR: operation has finished with an error

gaspi_atomic_compare_swap atomically compares the global value of the value on rank, segment segment_id_remote and offset offset_remote to the value of comparator. If the comparison is true, this global value is set to value_new. If the comparison is false, it keeps its value.

After successful procedure completion, i.e. return value GASPI_SUCCESS, the parameter value_old contains the previous value before the comparison was done.

If the procedure returns with GASPI_TIMEOUT, the compare and swap request could not be completed during the given timeout. The parameter value_old has an undefined value. A subsequent call of gaspi_atomic_compare_swap needs to be invoked in order to complete the operation.

If the procedure returns with GASPI_ERROR, the compare and swap request aborted abnormally. The parameter value_old as well as the global value (segment_id, offset, rank) have undefined values.

In both cases, GASPI_TIMEOUT and GASPI_ERROR, the GASPI state vector should be checked in order to deal with possible failures.

Implementor advice: The implementation might require some alignment restrictions that is, the triple(segment_id, offset, rank) might be required to respect some alignment restrictions.

User advice: Concurrent accesses to the location represented by the triple(segment_id, offset, rank) are possible but consistency must be handled by the application.

10.2.3 Examples

The example in listing 14 illustrates the usage of global atomic operations for implementing a global resource lock. The example is implemented with timeout.

Listing 14: Gaspi global resource lock implemented with atomic counters

```c
#include <GASPI.h>
#include <cassert.h>

#define SUCCESS_OR_RETURN(f) {
    const int ec = (f);
    if (ec != GASPI_SUCCESS) {
```

10.2 Atomic operation calls

```c
#define VAL_UNLOCKED 9999999

gaspi_return_t
global_lock_init ( const gaspi_segment_id_t seg,
                   const gaspi_offset_t off,
                   const gaspi_rank_t rank_loc,
                   , const gaspi_timeout_t timeout
)
{
    gaspi_rank_t iProc;
    SUCCESS_OR_RETURN (gaspi_proc_rank (&iProc));
    if( iProc == rank_loc)
    {
        gaspi_pointer_t vptr;
        gaspi_atomic_value_t *lock_ptr;
        SUCCESS_OR_RETURN(gaspi_segment_ptr, &vptr);
        lock_ptr = (gaspi_atomic_value_t *) vptr;
        *lock_ptr = VAL_UNLOCKED;
    }
    SUCCESS_OR_RETURN (gaspi_barrier ( GASPI_GROUP_ALL
                   , timeout
                   )
    );
    return GASPI_SUCCESS;
}

gaspi_return_t
global_try_lock ( const gaspi_segment_id_t seg,
                   const gaspi_offset_t off,
                   const gaspi_rank_t rank_loc,
                   , const gaspi_timeout_t timeout
)
{
    gaspi_rank_t iProc;
    SUCCESS_OR_RETURN (gaspi_proc_rank (&iProc));
    gaspi_atomic_value_t old_value;
    SUCCESS_OR_RETURN (gaspi_atomic_compare_swap ( seg
                   , &old_value,
                   &lock_ptr
                   )
    );
    return (old_value == VAL_UNLOCKED) ? GASPI_SUCCESS :
          GASPI_BUSY;
}
```

11 Collective communication

11.1 Introduction and overview

Collective operations are collective with respect to a given group. A necessary condition for successful collective procedure completion is that all GASPI processes forming the given group have invoked the operation.

11.1 Introduction and overview

Collective operations support both synchronous and asynchronous implementations as well as time-based blocking. That means, progress towards successful procedure completion can be achieved either inside the call (for a synchronous implementation) or outside of the call (for an asynchronous implementation) before the procedure exits. In the case of a timeout (which is indicated by return value GASPI_TIMEOUT) the operation is then continued in the next call of the procedure. This implies that a collective operation may involve several procedure calls until completion. Completion is indicated by return value GASPI_SUCCESS.

Collective operations are exclusive per group, i.e. only one collective operation of a specific type on a given group can run at a given time. Starting a specific collective operation before another one of the same kind is not finished on all processes of the group (and marked as such) is not allowed and yields undefined behavior. For example, two allreduce operations for one group can not run at the same time; however, an allreduce and a barrier operation can run at the same time.

The timeout is a necessary condition in order to be able to write failure tolerant code. Timeout = 0 makes an atomic portion of progress in the operation if possible. If progress is possible, the procedure returns as soon as the atomic portion of progress is achieved. Otherwise, the procedure returns immediately. Here, an atomic portion of progress is defined as the smallest set of non-dividable instructions in the current state of the collective operation.

Reduction operations can be defined by the application via callback functions.

User advice: Not every collective operation will be implementable in an asynchronous fashion – for example if a user-defined callback function is used within a global reduction. Progress in this case can only be achieved inside of the call. Especially for large systems this implies that a collective potentially has to be called a substantial number of times in order to complete – especially if used in combination with GASPI_TEST. In this combination the called collective immediately returns (after completing local work) and never waits for data from remote processes. A corresponding code fragment in this case would assume the form:

```c
while (GASPI_allreduce_user ( buffer_send, buffer_receive
, char num
, size_element
, reduce_operation
, reduce_state
, group
, GASPI_TEST
) != GASPI_SUCCESS)
{
    work();
}
```
11.2 Barrier synchronisation

11.2.1 gaspi_barrier

The gaspi_barrier procedure is a collective time-based blocking procedure. An implementation is free to provide it as a synchronous or an asynchronous procedure.

\[
\text{GASPI\_BARRIER ( group, timeout )}
\]

Parameter:

\( (\text{in}) \) group: the group of ranks which should participate in the barrier

\( (\text{in}) \) timeout: the timeout

\[
\text{gaspi\_return\_t}
\text{gaspi\_barrier ( gaspi\_group\_t group, gaspi\_timeout\_t timeout )}
\]

\[
\text{function gaspi\_barrier(group,timeout\_ms) \&}
\text{& result( res ) bind(C, name="gaspi\_barrier")}
\text{integer(gaspi\_group\_t), value :: group}
\text{integer(gaspi\_timeout\_t), value :: timeout\_ms}
\text{integer(gaspi\_return\_t) :: res}
\text{end function gaspi\_barrier}
\]

Execution phase:

Working

Return values:

GASPI\_SUCCESS: operation has returned successfully
GASPI\_TIMEOUT: operation has run into a timeout
GASPI\_ERROR: operation has finished with an error

\text{gaspi\_barrier} blocks the caller until all group members of group have invoked the procedure or if timeout milliseconds have been reached since procedure invocation. After successful procedure completion, i.e. return value GASPI\_SUCCESS, all group members have invoked the procedure. In case of GASPI\_TIMEOUT it is unknown whether or not all GASPI processes forming the given group have invoked the call.

Progress towards successful gaspi_barrier completion is achieved even if the procedure exits due to timeout. The barrier is then continued in the next call of the procedure. This implies that a barrier operation may involve several gaspi_barrier calls until completion.
Barrier operations are exclusive per group, i.e. only one barrier operation on a given group can run at a time. Starting a barrier operation in another thread before a previously invoked barrier is finished on all processes of the group is not allowed and yields undefined behavior.

In case of error, the return value is GASPI_ERROR. The error vector should be investigated.

User advice: The barrier is supposed to synchronise processes and not threads.

### 11.2.2 Examples

In the following example a gaspi_barrier is interrupted after 100 ms in order to check for errors.

```c
1  gaspi_return_t err;
2  do
3    { err = gaspi_barrier (g, 100);
4      if (err == GASPI_TIMEOUT && error vector indicates error)
5        { goto ERROR_HANDLING; }
6    } while (err != GASPI_SUCCESS);
```

The following example shows a non-blocking barrier. Some local work (in this case: cleanup) is performed, overlapping it with the barrier and only then a full synchronisation is achieved by calling the barrier again with a blocking semantics (if needed).

```c
1  const gaspi_return_t err = gaspi_barrier (g, GASPI_TEST);
2  do_local_cleanup();
3  if (err != GASPI_ERROR && err != GASPI_SUCCESS)
4    { gaspi_barrier (g, GASPI_BLOCK); }
```

### 11.3 Predefined global reduction operations

#### 11.3.1 gaspi_allreduce

The gaspi_allreduce procedure is a collective time-based blocking procedure. An implementation is free to provide it as a synchronous or an asynchronous procedure.
11.3 Predefined global reduction operations

**Parameter:**

(in) **buffer_send**: pointer to the buffer where the input is placed

(in) **buffer_receive**: pointer to the buffer where the result is placed

(in) **num**: the number of elements to be reduced on each process

(in) **operation**: the GASPI reduction operation type

(in) **datatype**: the GASPI element type

(in) **group**: the group of ranks which participate in the reduction operation

(in) **timeout**: the timeout

**Execution phase:**

**Return values:**

GASPI_SUCCESS: operation has returned successfully
11.3 Predefined global reduction operations

GASPI_TIMEOUT: operation has run into a timeout
GASPI_ERROR: operation has finished with an error

gaspi_allreduce combines the num elements of type datatypen residing in buffer_send on each process in accordance with the given operation. The reduction operation is on a per element basis, i.e. the operation is applied to each of the elements. gaspi_allreduce blocks the caller until all data is available that is needed to calculate the result or if timeout milliseconds have been reached since procedure invocation. After successful procedure completion, i.e. return value GASPI_SUCCESS, all group members have invoked the procedure and buffer_receive contains the result of the reduction operation on every GASPI process of group. In case of GASPI_TIMEOUT not all data is available that is needed to calculate the result.

Progress towards successful gaspi_allreduce completion is achieved even if the procedure exits due to timeout. The reduction operation is then continued in the next call of the procedure. This implies that a reduction operation may involve several gaspi_allreduce calls until completion.

Reduction operations are exclusive per group, i.e. only one reduction operation on a given group can run at a time. Starting a reduction operation for the same group in a separate thread before previously invoked operation is finished on all processes of the group is not allowed and yields undefined behavior.

The buffer_send as well as the buffer_receive do not need to reside in the global address space. gaspi_allreduce copies the send buffer into an internal buffer at the first invocation. The result is copied from an internal buffer into the receive buffer immediately before the procedure returns successfully. The buffers need to have the appropriate size to host all of the num elements. Otherwise the reduction operation yields undefined behavior. The maximum permissible number of elements is implementation dependent and can be retrieved by gaspi_allreduce_elem_max.

In case of error, the return value is GASPI_ERROR. The error vector should be examined. buffer_receive has an undefined value.

In case of GASPI_TIMEOUT, the reduction operation is not finished yet, i.e. not all data is available that is needed to calculate the result. The buffer_receive has an undefined value.

11.3.2 Predefined reduction operations

There are three predefined reduction operations:

```c
typedef enum { GASPI_OP_MIN
               , GASPI_OP_MAX
               , GASPI_OP_SUM
            } gaspi_operation_t;
```

GASPI_OP_MIN determines the minimum of the elements of each column of the input vector.
11.4 User-defined global reduction operations

**GASPI_OP_MAX** determines the maximum of the elements of each column of the input vector.

**GASPI_OP_SUM** sums up all elements of each column of the input vector.

### 11.3.3 Predefined types

And the types are:

```c
typedef enum { GASPI_TYPE_INT,
               GASPI_TYPE_UINT,
               GASPI_TYPE_LONG,
               GASPI_TYPE_ULONG,
               GASPI_TYPE_FLOAT,
               GASPI_TYPE_DOUBLE
} gaspi_datatype_t;
```

**GASPI_TYPE_INT** integer

**GASPI_TYPE_UINT** unsigned integer

**GASPI_TYPE_LONG** long

**GASPI_TYPE_ULONG** unsigned long

**GASPI_TYPE_FLOAT** float

**GASPI_TYPE_DOUBLE** double

11.4 User-defined global reduction operations

#### 11.4.1 gaspi_allreduce_user

The procedure `gaspi_allreduce_user` allows the user to specify its own reduction operation. Only operations are supported which are commutative and associative. It is a **collective time-based blocking** procedure. An implementation is free to provide it as a synchronous or an asynchronous procedure.

```c
GASPI_ALLREDUCE_USER ( buffer_send,
                        buffer_receive,
                        num,
                        size_element,
                        reduce_operation,
                        reduce_state,
                        group,
                        timeout )
```

**Parameter:**

*(in) buffer_send:* pointer to the buffer where the input is placed
11.4 User-defined global reduction operations

(in) buffer.receive: pointer to the buffer where the result is placed
(in) num: the number of elements to be reduced on each process
(in) size_element: Size in bytes of one element to be reduced
(in) reduce_operation: pointer to the user defined reduction operation procedure
(inout) reduce_state: reduction state vector
(in) group: the group of ranks which participate in the reduction operation
(in) timeout: the timeout

```
gaspi_return_t
gaspi_allreduce_user ( gaspi_const_pointer_t buffer_send
  , gaspi_pointer_t buffer_receive
  , gaspi_number_t num
  , gaspi_size_t size_element
  , gaspi_reduce_operation_t reduce_operation
  , gaspi_reduce_state_t reduce_state
  , gaspi_group_t group
  , gaspi_timeout_t timeout )
```

```
function gaspi_allreduce_user(buffer_send,buffer_receive, &
  & num,element_size,reduce_operation,reduce_state,&
  & group,timeout_ms) &
  & result( res ) bind(C, name="gaspi_allreduce_user")
type(c_ptr), value :: buffer_send
type(c_ptr), value :: buffer_receive
integer(gaspi_number_t), value :: num
integer(gaspi_size_t), value :: element_size
type(c_funptr), value :: reduce_operation
type(c_ptr), value :: reduce_state
integer(gaspi_group_t), value :: group
integer(gaspi_timeout_t), value :: timeout_ms
integer(gaspi_return_t) :: res
end function gaspi_allreduce_user
```

Execution phase:

Working

Return values:

GASPI_SUCCESS: operation has returned successfully
GASPI_TIMEOUT: operation has run into a timeout
GASPI_ERROR: operation has finished with an error

```
gaspi_allreduce_user has the same semantics as the predefined reduction
operation gaspi_allreduce described in the last section.
```
A user defined reduction operation *reduce_operation* and a user defined state *reduce_state* are passed.

The elements on which the user defined reduction operation is applied are described by their byte size *size_element*. The entire size of the data to be reduced, i.e. *num* times *size_element*, must not be larger than the internal buffer size of *gasi_allreduce_user*. The internal buffer size can be queried through *gasi_allreduce_buf_size*.

### 11.4.2 User defined reduction operations

The prototype for the user defined reduction operations is the following:

```c
GASPI_REDUCE_OPERATION ( operand_one , operand_two , result , state , timeout )
```

**Parameter:**

*(in) operand_one:* pointer to the first operand

*(in) operand_two:* pointer to the second operand

*(in) result:* pointer to the result

*(in) state:* pointer to the state

*(in) timeout:* the timeout

```c
gasi_return_t gaspi_reduce_operation ( gaspi_const_pointer_t operand_one , gaspi_const_pointer_t operand_two , gaspi_pointer_t result , gaspi_reduce_state_t state , gaspi_timeout_t timeout )
```
11.4 User-defined global reduction operations

function my_reduce_operation(op_one, op_two, op_res, &
& op_state, num, element_size, timeout) &
& result (res) bind(C, name="my_reduce_operation")
    implicit none
    integer(gaspi_number_t), intent(in), value :: num

! the fortran user defined callback function requires an explicit type from the iso_c_binding module. in this example integer(c_int) (op_one, op_two, op_res, op_state)

integer(c_int), intent(in) :: op_one(num)
integer(c_int), intent(in) :: op_two(num)
integer(c_int), intent(out) :: op_res(num)
integer(c_int), intent(out) :: op_state(num)
integer(gaspi_size_t), value :: element_size
integer(gaspi_timeout_t), value :: timeout
integer(gaspi_return_t) :: res

! your user defined operation
! ...
    res = GASPI_SUCCESS
end function my_reduce_operation

Return values:

GASPI_SUCCESS: operation has returned successfully
GASPI_TIMEOUT: operation has run into a timeout
GASPI_ERROR: operation has finished with an error

A pointer to the first operand and a pointer to the second operand are passed. The result is stored in the memory represented by the pointer result. In addition to the actual data, also a state can be passed to the operator which might be required in order to compute the result. In order to meet real time system specifications, a timeout can be passed to the user defined reduction operator. The reduction operator should return a gaspi_return_t with the same semantics, i.e. GASPI_SUCCESS for successful procedure completion, GASPI_TIMEOUT in case of timeout and GASPI_ERROR in case of error.

The user defined reduction operator needs to be commutative and associative. The reduce operator type passed to gaspi_allreduce_user is a pointer to a function with the prototype described above.

typedef gaspi_reduce_operation* gaspi_reduce_operation_t

The GASPI reduction operation type

11.4.3 allreduce state

The allreduce state type
typedef void* gaspi_reduce_state_t

The GASPI reduction operation state type

is a pointer to a state which may be passed to the user defined reduction operation. A state may contain additional information beside the actual data to be reduced needed to perform the reduction operation.

11.4.4 Example

A fortran version of the user defined allreduce hence might assume the form listing 15

Listing 15: GASPI User defined allreduce, fortran example.

```fortran
module my_reduce

  use gaspi_c_binding
  implicit none
  contains

  function my_reduce_operation(op_one, op_two, op_res, &
    op_state, num, element_size, timeout) &
    result (res) bind(C,name="my_reduce_operation")
    implicit none
    integer(gaspi_number_t), intent(in), value :: num
    integer(c_int), intent(in) :: op_one(num)
    integer(c_int), intent(in) :: op_two(num)
    integer(c_int), intent(out) :: op_res(num)
    integer(c_int), intent(out) :: op_state(num)
    integer(gaspi_size_t), value :: element_size
    integer(gaspi_timeout_t), value :: timeout
    integer(gaspi_return_t) :: res

    integer i
    do i = 1, num
      op_res(i) = max(op_one(i), op_two(i))
    enddo
    res = GASPI_SUCCESS
  end function my_reduce_operation

end module my_reduce

program allreduce

  use gaspi_c_binding
  use my_reduce
  implicit none
  integer(gaspi_size_t) :: sizeof_int
  integer(gaspi_return_t) :: res
```
The **Gaspi** specification provides getter functions for all entries in the **Gaspi** configuration. These getter functions are *synchronous local blocking* procedures which, after successful procedure completion (i.e., return value **GASPI_SUCCESS**), read out the corresponding value of the current configuration setting.

The values of the parameters in the **Gaspi** configuration are determined in **gaspi_proc_init** at startup. If the value of one of these parameters is compliant with the system capabilities, the parameter is set to the requested/preferred value. Otherwise, the parameter is set to the maximum value compliant with the system capabilities. The values of the parameters realised in the **Gaspi** configuration are implementation specific.

In case of error, the return value is **GASPI_ERROR** and the corresponding parameter in the getter function has an undefined value.
12.1 Getter functions for group management

12.1.1 gaspi_group_max

GASPI_GROUP_MAX (group_max)

Parameter:

(out) group_max: the total number of groups

gaspi_return_t

gaspi_group_max (gaspi_number_t *group_max)

function gaspi_group_max(group_max) &
& result ( res ) bind(C, name="gaspi_group_max")
  integer(gaspi_number_t) :: group_max
  integer(gaspi_return_t) :: res
end function gaspi_group_max

Execution phase:

Working

Return values:

GASPI_SUCCESS: operation has returned successfully
GASPI_ERROR: operation has finished with an error

12.2 Getter functions for segment management

12.2.1 gaspi_segment_max

GASPI_SEGMENT_MAX (segment_max)

Parameter:

(out) segment_max: the total number of permissible segments

gaspi_return_t

gaspi_segment_max (gaspi_number_t *segment_max)

function gaspi_segment_max(segment_max) &
& result ( res ) bind(C, name="gaspi_segment_max")
  integer(gaspi_number_t) :: segment_max
  integer(gaspi_return_t) :: res
end function gaspi_segment_max

Execution phase:
12.3 Getter functions for communication management

12.3.1 gaspi_queue_num

\[ \text{GASPI\_QUEUE\_NUM} \ (\text{queue\_num}) \]

**Parameter:**

\( \text{(out) queue\_num: the number of available queues} \)

\[ \text{gaspi\_return\_t gaspi\_queue\_num} \ (\text{gaspi\_number\_t} \ast \text{queue\_num}) \]

\[ \text{function gaspi\_queue\_num(queue\_num) \&} \]
\[ \text{\quad result( res ) bind(C, name=\text{"gaspi\_queue\_num"})} \]
\[ \text{\quad integer(gaspi\_number\_t) :: queue\_num} \]
\[ \text{\quad integer(gaspi\_return\_t) :: res} \]
\[ \text{end function gaspi\_queue\_num} \]

**Execution phase:**

Working

**Return values:**

GASPI\_SUCCESS: operation has returned successfully
GASPI\_ERROR: operation has finished with an error

12.3.2 gaspi_queue_size_max

\[ \text{GASPI\_QUEUE\_SIZE\_MAX} \ (\text{queue\_size\_max}) \]

**Parameter:**

\( \text{(out) queue\_size\_max: the maximum number of simultaneous requests allowed} \)

\[ \text{gaspi\_return\_t gaspi\_queue\_size\_max} \ (\text{gaspi\_number\_t} \ast \text{queue\_size\_max}) \]

\[ \text{function gaspi\_queue\_size\_max(queue\_size\_max) \&} \]
\[ \text{\quad result( res ) bind(C, name=\text{"gaspi\_queue\_size\_max"})} \]
\[ \text{\quad integer(gaspi\_number\_t) :: queue\_size\_max} \]
\[ \text{end function gaspi\_queue\_size\_max} \]
12.3 Getter functions for communication management

`function gaspi_queue_size_max(queue_size_max) &
  &     result( res ) bind(C, name="gaspi_queue_size_max")
  integer(gaspi_number_t) :: queue_size_max
  integer(gaspi_return_t) :: res
end function gaspi_queue_size_max`

*Execution phase:*
Working

*Return values:*
GASPI_SUCCESS: operation has returned successfully
GASPI_ERROR: operation has finished with an error

12.3.3 gaspi_transfer_size_max

`GASPI_TRANSFER_SIZE_MAX (transfer_size_max)`

*Parameter:*
*(out)* `transfer_size_max`: the maximum transfer size allowed for a single request.

`gaspi_return_t
  gaspi_transfer_size_max (gaspi_size_t *transfer_size_max)`

`function gaspi_transfer_size_max(transfer_size_max) &
   &     result( res ) &
   &     bind(C, name="gaspi_transfer_size_max")
   integer(gaspi_size_t) :: transfer_size_max
   integer(gaspi_return_t) :: res
end function gaspi_transfer_size_max`

*Execution phase:*
Working

*Return values:*
GASPI_SUCCESS: operation has returned successfully
GASPI_ERROR: operation has finished with an error

12.3.4 gaspi_notification_num

`GASPI_NOTIFICATION_NUM (notification_num)`

*Parameter:*


12.4 Getter functions for passive communication

(out) \textit{notification\_num}: the number of available notifications

\begin{verbatim}
gaspi\_return\_t
gaspi\_notification\_num (gaspi\_number\_t \ast notification\_num)

function gaspi\_notification\_num(notification\_num) &
& result( res ) bind(C, name="gaspi\_notification\_num")
  integer(gaspi\_number\_t) :: notification\_num
  integer(gaspi\_return\_t) :: res
end function gaspi\_notification\_num
\end{verbatim}

\textit{Execution phase:}

\begin{itemize}
  \item Working
\end{itemize}

\textit{Return values:}

\begin{itemize}
  \item GASPI\_SUCCESS: operation has returned successfully
  \item GASPI\_ERROR: operation has finished with an error
\end{itemize}

12.4 Getter functions for passive communication

12.4.1 gaspi\_passive\_transfer\_size\_max

\begin{verbatim}
GASPI\_PASSIVE\_TRANSFER\_SIZE\_MAX (transfer\_size\_max)

Parameter:

(out) \textit{transfer\_size\_max}: maximal transfer size per single passive communication request

\begin{verbatim}
gaspi\_return\_t
gaspi\_passive\_transfer\_size\_max (gaspi\_size\_t \ast transfer\_size\_max)

function gaspi\_passive\_transfer\_size\_max(transfer\_size\_max) &
& result( res ) &
& bind(C, name="gaspi\_passive\_transfer\_size\_max")
  integer(gaspi\_size\_t) :: transfer\_size\_max
  integer(gaspi\_return\_t) :: res
end function gaspi\_passive\_transfer\_size\_max
\end{verbatim}

\textit{Execution phase:}

\begin{itemize}
  \item Working
\end{itemize}

\textit{Return values:}

\begin{itemize}
  \item GASPI\_SUCCESS: operation has returned successfully
  \item GASPI\_ERROR: operation has finished with an error
\end{itemize}
12.5 Getter functions related to atomic operations

12.5.1 gaspi_atomic_max

```fortran
GASPI_ATOMIC_MAX (max_value)
```

Parameter:

(out) max_value: the maximum value an gaspi_atomic_value_t can hold

```fortran
gaspi_return_t
gaspi_atomic_max (gaspi_atomic_value_t *max_value)
```

```fortran
function gaspi_atomic_max(max_value) &
  &   result( res ) bind(C, name="gaspi_atomic_max")
  integer(gaspi_atomic_value_t) :: max_value
  integer(gaspi_return_t) :: res
end function gaspi_atomic_max
```

Execution phase:

Working

Return values:

GASPI_SUCCESS: operation has returned successfully
GASPI_ERROR: operation has finished with an error

12.6 Getter functions for collective communication

12.6.1 gaspi_allreduce_buf_size

```fortran
GASPI_ALLREDUCE_BUF_SIZE (buf_size)
```

Parameter:

(out) buf_size: the size of the internal buffer in gaspi_allreduce_user

```fortran
gaspi_return_t
gaspi_allreduce_buf_size (gaspi_size_t *buf_size)
```

```fortran
function gaspi_allreduce_buf_size(buf_size) &
  &   result( res ) bind(C, name="gaspi_allreduce_buf_size")
  integer(gaspi_size_t) :: buf_size
  integer(gaspi_return_t) :: res
end function gaspi_allreduce_buf_size
```

Execution phase:
12.7 Getter functions related to infrastructure

Working

Return values:
GASPI_SUCCESS: operation has returned successfully
GASPI_ERROR: operation has finished with an error

12.6.2 gaspi_allreduce_elem_max

GASPI_ALLREDUCE_ELEM_MAX (elem_max)

Parameter:
(out) elem_max: the maximum number of elements allowed in gaspi_allreduce

```
gaspi_return_t
gaspi_allreduce_elem_max (gaspi_number_t *elem_max)
```

function gaspi_allreduce_elem_max(elem_max) &
&
       result ( res ) bind(C, name="gaspi_allreduce_elem_max")
       integer(gaspi_number_t) :: elem_max
       integer(gaspi_return_t) :: res
end function gaspi_allreduce_elem_max

Execution phase:

Working

Return values:
GASPI_SUCCESS: operation has returned successfully
GASPI_ERROR: operation has finished with an error

12.7 Getter functions related to infrastructure

12.7.1 gaspi_network_type

GASPI_NETWORK_TYPE (network_type)

Parameter:
(out) network_type: the chosen network type

```
gaspi_return_t
gaspi_network_type (gaspi_network_t *network_type)
```
12.7 Getter functions related to infrastructure

```
function gaspi_network_type(network_type) &
  &    result ( res ) bind(C, name="gaspi_network_type")
  integer(gaspi_network_t) :: network_type
  integer(gaspi_return_t) :: res
end function gaspi_network_type
```

**Execution phase:**
Working

**Return values:**
- GASPI_SUCCESS: operation has returned successfully
- GASPI_ERROR: operation has finished with an error

### 12.7.2 gaspi_build_infrastructure

```
GASPI_BUILD_INFRASTRUCTURE (build_infrastructure)
```

**Parameter:**
- (out) `build_infrastructure`: the current value of `build_infrastructure`

```
gaspi_return_t
gaspi_build_infrastructure (gaspi_number_t *build_infrastructure)
```

function gaspi_build_infrastructure(build_infrastructure) &
  &    result ( res ) &
  &    bind(C, name="gaspi_build_infrastructure")
  integer (gaspi_number_t) :: build_infrastructure
  integer(gaspi_return_t) :: res
end function gaspi_build_infrastructure

**Execution phase:**
Working

**Return values:**
- GASPI_SUCCESS: operation has returned successfully
- GASPI_ERROR: operation has finished with an error
13 GASPI Environmental Management

13.1 Implementation Information

13.1.1 gaspi_version

The `gaspi_version` procedure is a *synchronous local blocking* procedure which determines the version of the running GASPI installation.

```fortran
GASPI_VERSION (version)
```

*Parameter:*

*(out)* `version`: The version of the running GASPI installation

```fortran
gaspi_return_t
    gaspi_version (float *version)
```

```fortran
function gaspi_version(version) &
    result(res) bind(C, name="gaspi_version")
    real(gaspi_float) :: version
    integer(gaspi_return_t) :: res
end function gaspi_version
```

*Execution phase:*

Any

*Return values:*

GASPI_SUCCESS: operation has returned successfully
GASPI_ERROR: operation has finished with an error

After successful procedure completion, i.e. return value `GASPI_SUCCESS version` contains the version of the running GASPI installation.

In case of error, the return value is GASPI_ERROR. The output parameter `version` has an undefined value.

13.2 Timing information

13.2.1 gaspi_time_get

The `gaspi_time_get` procedure is a *synchronous local blocking* procedure which determines the time elapsed since an arbitrary point of time in the past.

```fortran
GASPI_TIME_GET (stime)
```
13.2 Timing information

Parameter:

*(out)* `wtime`: time elapsed in milliseconds

```
gaspi_return_t
gaspi_time_get (gaspi_time_t *wtime)
```

```
function gaspi_time_get(vtime) &
  & result( res ) bind(C, name="gaspi_time_get")
  integer(gaspi_time_t) :: vtime
  integer(gaspi_return_t) :: res
end function gaspi_time_get
```

Execution phase:

Working

Return values:

`GASPI_SUCCESS`: operation has returned successfully

`GASPI_ERROR`: operation has finished with an error

After successful procedure completion, i.e. return value `GASPI_SUCCESS`, the parameter `wtime` contains elapsed time in milliseconds since an arbitrary point in the past. The parameter `wtime` is not synchronised among the different Gaspi processes.

In case of error, the return value is `GASPI_ERROR`. The value of the output parameter `wtime` is undefined.

13.2.2 `gaspi_time_ticks`

The `gaspi_time_ticks` procedure is a *synchronous local blocking* procedure which returns the resolution of the internal timer in terms of milliseconds.

```
GASPI_TIME_TICKS (resolution)
```

Parameter:

*(out)* `resolution`: the resolution of the internal timer in milliseconds

```
gaspi_return_t
gaspi_time_ticks (gaspi_time_t *resolution)
```

```
function gaspi_time_ticks(resolution) &
  & result( res ) bind(C, name="gaspi_time_ticks")
  integer(gaspi_time_t) :: resolution
  integer(gaspi_return_t) :: res
end function gaspi_time_ticks`
13.3 Error Codes and Classes

**Execution phase:**
Any

**Return values:**
- GASPI_SUCCESS: operation has returned successfully
- GASPI_ERROR: operation has finished with an error

After successful procedure completion, i.e., return value GASPI_SUCCESS, the parameter *resolution* contains the resolution of the internal timer in units of milliseconds.

In case of error, the return value is GASPI_ERROR. The value of the output parameter *resolution* is undefined.

### 13.3 Error Codes and Classes

#### 13.3.1 GASPI error codes

In principle all return values less than zero represent an error. Every implementation is free to define specific error codes.

#### 13.3.2 gaspi_print_error

The *gaspi_print_error* procedure is a *synchronous local blocking* procedure which translates an error code to a text message.

```c
GASPI_PRINT_ERROR( error_code, error_message )
```

**Parameter:**

*(in) error_code:* the error code to be translated

*(out) error_message:* the error message

```c
function gaspi_print_error(error_code, error_message) &
& result( res ) bind(C, name="gaspi_print_error")
  integer(gaspi_return_t), value :: error_code
  character(c_char), dimension(*) :: error_message
  integer(gaspi_return_t) :: res
end function gaspi_print_error
```

**Execution phase:**
Any

Return values:
GASPI_SUCCESS: operation has returned successfully
GASPI_ERROR: operation has finished with an error

After successful procedure completion, i.e. return value GASPI_SUCCESS error_message contains the error message corresponding to the error code error_code.

In case of error, the return value is GASPI_ERROR.

The procedure can be invoked in any of the GASPI execution phases.

14 Profiling Interface

The profiling interface of GASPI consists of two parts. The statistics part provides the means to allow the user to collect basic profiling data about a program run. The event tracing part describes the requirements for an GASPI implementation in order to support the transparent interception and inspection of function calls.

14.1 Statistics

14.1.1 gaspi_statistic_counter_max

The gaspi_statistic_counter_max procedure is a synchronous local blocking procedure, which provides a way to inform the GASPI user dynamically about the number of available counters. An implementation should not provide a compile-time constant maximum for gaspi_statistic_counter_t. Instead the user can call gaspi_statistic_counter_max in order to determine the maximum value for gaspi_statistic_counter_t.

GASPI_STATISTIC_COUNTER_MAX ( counter_max )

Parameter:
(out) counter_max: the maximum value for gaspi_statistic_counter_t. The allowed value range is 0 ≤ counter < counter_max

gaspi_return_t
gaspi_statistic_counter_max ( gaspi_number_t *counter_max )
14.1 Statistics

function gaspi_statistic_counter_max(counter_max) &
   & result( res ) &
   & bind(C, name="gaspi_statistic_counter_max")
   integer(gaspi_statistic_counter_t) :: counter_max
   integer(gaspi_return_t) :: res
end function gaspi_statistic_counter_max

Execution phase:

Any

Return values:

GASPI_SUCCESS: operation has returned successfully
GASPI_ERROR: operation has finished with an error

If a GASPI implementation defines symbolic constants for gaspi_statistic_counter_t a priori, then gaspi_statistic_counter_max should set counter_max to the corresponding maximum value. A high-speed implementation will likely set counter_max to 0 and does not provide any statistics by default. A dynamically linked wrapper library can provide extra counters by adjusting the return value of gaspi_statistic_counter_max.

Library implementor advice: A sensible wrapper library will respect the value returned by the native gaspi_statistic_counter_max and append their own counters accordingly. Thus accesses to statistic counters provided by the GASPI implementation itself are not harmed.

14.1.2 gaspi_statistic_counter_info

The gaspi_statistic_counter_info procedure is a synchronous local blocking procedure which provides an implementation independent way to retrieve information for a particular statistic counter. Beside the name and a description this function also yields the meaning of the argument value for this counter, if any. The meaning is defined in terms of the gaspi_statistic_argument_t enumeration.

typedef enum { GASPI_STATISTIC_ARGUMENT_NONE
   , GASPI_STATISTIC_ARGUMENT_RANK
   , ...
} gaspi_statistic_argument_t;

A GASPI implementation is free to extend the above enumeration.

GASPI_STATISTIC_COUNTER_INFO ( const counter
   , argument
   , counter_name
   , counter_description
   , verbosity_level )
Parameter:

*(in) counter*: the counter, for which detailed information is requested

*(out) counter_argument*: the meaning of the argument value

*(out) counter_name*: a short name of this counter

*(out) counter_description*: a more verbose description of this counter

*(out) verbosity_level*: minimum verbosity level to activate this counter (at least 1)

```plaintext
function gaspi_statistic_counter_info(counter, counter_argument, &
 & counter_name, counter_description, verbosity_level) &
 & result(res) &
 & bind(C, name="gaspi_statistic_counter_info")
 integer(gaspi_statistic_counter_t), value:: counter
 integer(gaspi_statistic_argument_t):: counter_argument
 character(c_char), dimension(*):: counter_name
 character(c_char), dimension(*):: counter_description
 integer(gaspi_number_t):: verbosity_level
 integer(gaspi_return_t):: res
end function gaspi_statistic_counter_info
```

Execution phase:

Any

Return values:

GASPI_SUCCESS: operation has returned successfully

GASPI_ERROR: operation has finished with an error

After successful procedure completion, i.e. return value GASPI_SUCCESS, the output variables contain the desired information. A dynamically linked wrapper library should provide information for added counters by wrapping gaspi_statistic_counter_info. The verbosity level for all counters should be at least 1 (see gaspi_statistic_verbosity_level below).

If the return value is GASPI_ERROR, the particular counter issued to gaspi_statistic_counter_info does not exist.
14.1.3 gaspi_statistic_verbosity_level

The gaspi_statistic_verbosity_level procedure is a synchronous local blocking procedure which sets the process-wide verbosity level of the statistic interface. A counter is only active (that is, it is updated), if the process-wide verbosity level is higher or equal to the minimum verbosity level of that counter. If a call to gaspi_statistic_verbosity_level activates or deactivates counters and there are asynchronous operations in progress, it is unspecified, whether and how these counters are affected by the operations. It is furthermore unspecified whether and how counters of higher verbosity levels are updated.

A verbosity level of 0 deactivates all counting. It is not guaranteed, that counters with a minimum verbosity level of 0 are counted properly, if the verbosity level is set to 0.

GASPI_STATISTIC_VERBOSITY_LEVEL ( verbosity_level )

Parameter:

(in) verbosity_level: the level of desired verbosity

gaspi_statistic_verbosity_level ( gaspi_number_t verbosity_level)

function gaspi_statistic_verbosity_level(verbosity_level_) &
&       result( res ) &
&       bind(C, name="gaspi_statistic_verbosity_level")
       integer(gaspi_number_t), value :: verbosity_level_
       integer(gaspi_return_t) :: res
end function gaspi_statistic_verbosity_level

Execution phase:

Any

Return values:

GASPI_SUCCESS: operation has returned successfully
GASPI_ERROR: operation has finished with an error

14.1.4 gaspi_statistic_counter_get

The gaspi_statistic_counter_get procedure is a synchronous local blocking procedure which retrieves a statistical counter from the local GASPI process.

GASPI_STATISTIC_COUNTER_GET ( counter
   , argument
   , value )
14.1 Statistics

Parameter:

(in) counter: the counter to be retrieved

(in) argument: the argument for the counter

(out) value: the current value of the counter

```fortran
function gaspi_statistic_counter_get(counter,argument,&
   value_arg) &
   result(res) &
   bind(C, name="gaspi_statistic_counter_get")
   integer(gaspi_statistic_counter_t), value :: counter
   integer(gaspi_statistic_argument_t), value :: argument
   integer(gaspi_number_t) :: value_arg
   integer(gaspi_return_t) :: res
end function gaspi_statistic_counter_get
```

Execution phase:

Any

Return values:

GASPI_SUCCESS: operation has returned successfully

GASPI_ERROR: operation has finished with an error

The meaning of parameter argument depends on the retrieved counter. For instance, if a counter retrieves the bytes sent per target rank, then argument contains the target rank number. If the retrieved counter has no argument, the value of argument is ignored. After successful procedure completion, i.e. return value GASPI_SUCCESS value contains the current value of the corresponding counter.

The return value is GASPI_ERROR, if counter does not exist, i.e. exceeds gaspi_statistic_counter_max.

It is allowed to access a counter even, if the process-wide verbosity level is lower than the minimum verbosity level of that counter. Thus it is possible to profile certain regions of an application by changing the verbosity level and read the counter values at a later point in time independently of the current verbosity level.

14.1.5 gaspi_statistic_counter_reset

The gaspi_statistic_counter_reset procedure is a synchronous local blocking procedure which sets a statistical counter to 0.
14.2 Event Tracing

The **Gaspi** event tracing interface defines the requirements for an implementation to support the transparent interception and inspection of **Gaspi** calls. A **Gaspi** implementation must provide a mechanism, through which all **Gaspi** functions may be accessed with a name shift. The alternate entry point names have the prefix `pgaspi_` instead of `gaspi_`. In addition, the function `gaspi_pcontrol` is provided.

### 14.2.1 `gaspi_pcontrol`

The function `gaspi_pcontrol` is a no-op. A **Gaspi** implementation itself ignores the value of `argument` and returns immediately.

This routine is provided in order to enable users to communicate with an event trace interface from inside the application. The meaning of `argument` is specified by the used event tracer.

```fortran
GASPI_PCONTROL ( argument )
```

**Parameter:**

*(in) counter:* the counter to be reset

```fortran
function gaspi_statistic_counter_reset(counter) &
&   result(res) &
&   bind(C, name="gaspi_statistic_counter_reset")
   integer(gaspi_statistic_counter_t), value :: counter
   integer(gaspi_return_t) :: res
end function gaspi_statistic_counter_reset
```

**Execution phase:**

Any

**Return values:**

GASPI_SUCCESS: operation has returned successfully

GASPI_ERROR: operation has finished with an error

The return value is GASPI_ERROR, if `counter` does not exist, i.e. exceeds `gaspi_statistic_counter_max`.

---

**GASPI_STATISTIC_COUNTER_RESET (counter)**

**Parameter:**

*(in) counter:* the counter to be reset

```
GASPI_STATISTIC_COUNTER_RESET (counter)
```
Parameter:

(inout) argument:

gaspi_return_t gaspi_pcontrol ( gaspi_pointer_t argument )

function gaspi_pcontrol(argument) &
& result(res) bind(C, name="gaspi_pcontrol")
  type(c_ptr), value:: argument
  integer(gaspi_return_t):: res
end function gaspi_pcontrol

Execution phase:

Any

Return values:

GASPI_SUCCESS: operation has returned successfully
GASPI_ERROR: operation has finished with an error

A Listings

A.1 success_or_die

Listing 16: success_or_die.h

```c
#ifndef _SUCCESS_OR_DIE_H
#define _SUCCESS_OR_DIE_H

void success_or_die ( const char* file, const int line, const int ec );

#endif
```

Listing 17: success_or_die.c

```c
#include <success_or_die.h>
#include <stdlib.h>
#include <stdio.h>
```
### A.2 wait_if_queue_full

Listing 18: wait_if_queue_full.h

```c
#include <GASPI.h>

void success_or_die (const char* file, const int line, const int ec) {
    if (ec != GASPI_SUCCESS) {
        gaspi_string_t str;
        gaspi_error_message(ec, &str);
        fprintf(stderr, "error in %s[%i]: %s\n", file, line, str);
        exit(EXIT_FAILURE);
    }
}
```

Listing 19: wait_if_queue_full.c

```c
#include <wait_if_queue_full.h>
#include <success_or_die.h>

void wait_if_queue_full (const gaspi_queue_id_t queue_id, const gaspi_number_t request_size) {
    gaspi_number_t queue_size_max;
    gaspi_number_t queue_size;
    ASSERT(gaspi_queue_size_max (&queue_size_max));
    ASSERT(gaspi_queue_size (queue_id, &queue_size));
    if (queue_size + request_size >= queue_size_max)
```
{ 
    ASSERT (gaspi_wait (queue_id, GASPI_BLOCK));
}